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Multiobjective Metaheuristic Approaches
for Mean-Risk Combinatorial Optimisation
with Applications to Capacity Expansion

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à Teresa
ao Daniel
ao Pedro
à Ana
à Maria

Resumo

Muitas decisões em Gestão de Operações, em particular a um nível estratégico, são tomadas na presença de incerteza. Tendo em conta o impacto destas decisões, a questão do risco está surpreendentemente ausente da maioria da investigação e do trabalho aplicado nesta área. Tal poderá ser parcialmente explicado pela complexidade dos modelos de optimização para estes problemas, uma vez que necessitam de incluir parâmetros incertos, variáveis de decisão lógicas ou outras de natureza discreta, e mais do que um objectivo.

Uma das áreas de decisão críticas no âmbito da Estratégia de Operações é a área da Expansão de Capacidade, que se ocupa das decisões quanto ao tipo, dimensão, calendarização e localização dos investimentos em capacidade. Os modelos de capacidade têm de tratar diversas questões relacionadas com a complexidade acima referida, questões estas que conduzem a não-linearidades, não-convexidades, integridade e objectivos múltiplos.

Por outro lado, as meta-heurísticas multiobjectivo são algoritmos de optimização com características que favorecem uma aplicação extremamente eficiente em problemas com estas dificuldades, tendo, por este motivo, o potencial de vir a assumir um papel importante como abordagens genéricas para problemas de optimização combinatoria simultaneamente envolvendo a optimização do valor médio dos resultados e a minimização do risco. O principal objectivo deste trabalho foi realizar uma avaliação preliminar deste potencial.

Na primeira parte desta dissertação, apresenta-se uma abordagem de meta-heurísticas multiobjectivo baseadas em pesquisa local para uma *formulação média-risco de um problema da mochila estocástico estático*, considerando uma versão exacta e uma versão com aproximação amostral do problema, e a variância e o valor em risco condicional como medidas de risco.

A segunda parte deste trabalho debruça-se sobre uma *formulação média-risco para um problema de investimento em capacidade multi-período*, com irreversibilidade, indivisibilidade e economias de escala nos custos de capacidade. É, para este problema, proposta uma abordagem de meta-heurísticas multiobjectivo baseadas em pesquisa local, considerando o valor em risco condicional como medida de risco.

Na terceira parte da dissertação, introduz-se *flexibilidade de processo* no problema tratado na segunda parte, o que conduz, em cada período, a decisões de natureza discreta relativas ao investimento em expansão de capacidade, e decisões de natureza contínua relativas à utilização da capacidade disponível para satisfazer a procura. Os problemas de utilização de capacidade são resolvidos com programação linear, com o objectivo de determinar a capacidade mínima exigida para cada recurso, quando os restantes permanecem inalterados, disponibilizando, assim, informação sobre a admissibilidade das decisões de investimento. A este problema são aplicadas meta-heurísticas multiobjectivo baseadas em pesquisa local (em que novamente se considera o valor em risco condicional como medida de risco).

Os estudos computacionais realizados indicam claramente que as abordagens desenvolvidas são capazes de produzir aproximações aos conjuntos eficientes média-risco de elevada qualidade, com um esforço computacional modesto. Fica, assim, validada a hipótese de que as meta-heurísticas multiobjectivo constituem uma classe de algoritmos apropriados para lidar com as dificuldades apresentadas pelos problemas de optimização combinatoria média-risco.

Abstract

Many decisions in Operations Management, in particular at a strategic level, are made in the presence of uncertainty. Considering the impact of these decisions, risk concerns are surprisingly absent in the majority of research and applied work in this area. This may be partially explained by the complexity of optimisation models for these problems, as they must include uncertain parameters, logical or other discrete decision variables, and more than one objective.

One of the critical decision areas within Operations Strategy is Capacity Expansion, which is concerned with deciding the type, magnitude, timing, and location of capacity acquisition. Capacity models are required to address a variety of problem features related to the previously mentioned complexity, these features leading to nonlinearities, nonconvexities, integrality and multiple objectives.

On the other hand, multiobjective metaheuristics are optimisation algorithms extremely well suited to efficiently tackle problems that present these difficulties. They have therefore the potential to play an important role as general approaches for combinatorial optimisation problems simultaneously dealing with the optimisation of mean results and the minimisation of risk. The primary objective of our work was to perform a preliminary assessment of this potential.

In the first part of this dissertation, we present a multiobjective local search metaheuristic approach for both exact and sample approximation versions of a *mean-risk static stochastic knapsack problem*, considering both variance and conditional

value-at-risk as risk measures.

The second part of this work is concerned with a *mean-risk multistage capacity investment problem* with irreversibility, lumpiness and economies of scale in capacity costs. We propose a multiobjective local search metaheuristic approach for this problem, considering conditional value-at-risk as a risk measure.

In the third part of the dissertation, we introduce *process flexibility* in the problem addressed in the second part, leading to the consideration, in each period, of discrete decisions concerning the investment in capacity expansion, and continuous decisions concerning the utilization of the available capacity to satisfy demand. We solve the capacity utilization problems with linear programming, in order to find the minimum capacity for each resource with the other resources remaining unchanged. In this way, information is provided on the feasibility of the discrete investment decisions. We apply a multiobjective local search metaheuristic to this problem, again considering conditional value-at-risk as a risk measure.

Results of computational studies are presented, that clearly indicate the designed approaches are capable of producing high-quality approximations to the mean-risk efficient sets, with a modest computational effort, thus validating the hypothesis that multiobjective metaheuristics are a class of algorithms well suited to deal with the difficulties presented by mean-risk combinatorial optimisation problems.

Résumé

Beaucoup de décisions en Gestion des Opérations, en particulier au niveau stratégique, sont prises en présence d'incertitude. En considérant l'impact de ces décisions, c'est surprenant que la question du risque soit absente de la majorité de la recherche et du travail appliqué dans ce domaine. Ceci peut être partiellement expliqué par la complexité des modèles d'optimisation pour ces problèmes, qui demandent l'inclusion de paramètres incertains, variables de décision logiques ou autres de nature discrète, et plusieurs objectifs.

Un des principaux secteurs de décision en Stratégie des Opérations, c'est celui de l'Expansion de Capacité, qui s'occupe des décisions sur le type, la dimension, le calendrier et la localisation des investissements en capacité. Les modèles de capacité doivent considérer plusieurs aspects liés à la complexité mentionnée, aspects qui conduisent à l'existence de non-linéarités et de non-convexités, à l'intégralité des variables et à des objectifs multiples.

D'autre part, les métaheuristiques multiobjectif sont des algorithmes d'optimisation très bien adaptés à la résolution efficace des problèmes avec ces difficultés, et elles ont, pour cette raison, le potentiel de jouer un rôle important comme approches génériques pour des problèmes d'optimisation combinatoire traitant simultanément l'optimisation de la moyenne des résultats et la minimisation du risque. Le principal objectif de notre travail a été de réaliser une évaluation préliminaire de ce potentiel.

Dans la première partie de ce travail, nous présentons une approche de méta-

heuristiques multiobjectif basées sur recherche locale pour une *formulation moyenne-risque d'un problème de sac à dos stochastique statique*, en considérant une version exacte et une version avec approximation par échantillonnage, et la variance et la valeur à risque conditionnelle comme mesures de risque.

La deuxième partie de ce travail aborde une *formulation moyenne-risque pour un problème d'investissement en capacité multipériode*, avec irréversibilité, indivisibilité et économies d'échelle dans les coûts de capacité. Pour ce problème, nous proposons une approche de métaheuristiques multiobjectif basées sur recherche locale, en considérant la valeur à risque conditionnelle comme mesure de risque.

Dans la troisième partie de ce travail, nous ajoutons *flexibilité de processus* au problème abordé dans la deuxième partie, ce qui nous mène à considérer, dans chaque période, des décisions de nature discrète associées à l'investissement en expansion de capacité, et des décisions de nature continue associées à l'utilisation de la capacité disponible pour satisfaire la demande. Nous résolvons les problèmes d'utilisation de capacité par programmation linéaire, pour trouver le minimum de capacité requis pour chaque ressource, tandis que les autres ressources sont fixées. Cette procédure fournit de l'information sur l'admissibilité des décisions d'investissement. Nous appliquons à ce problème des métaheuristiques multiobjectif basées sur recherche locale, en considérant encore la valeur à risque conditionnelle comme mesure de risque.

Les études computationnelles réalisées indiquent clairement que les approches développées sont capables de produire des approximations aux ensembles efficaces moyenne-risque de qualité, avec un petit effort computationnel, en validant l'hypothèse de que les métaheuristiques multiobjectif sont des algorithmes d'optimisation très bien adaptés pour traiter les difficultés présentées par les problèmes d'optimisation combinatoire moyenne-risque.

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*Quisiera a veces borrar todos mis versos
para escribir por primera vez un poema.*

*Todo lo escrito no me alcanza
para sentir que he escrito uno.*

*Tampoco es suficiente haber vivido:
vivir comienza siempre ahora.*

Roberto Juarroz

Decimotercera Poesía Vertical

Chapter 1

Introduction

1.1 Research Opportunity

A large number of decisions in Operations Management are made in the presence of uncertainty. In fact, key factors, such as prices, resource availability or product demand, are regularly characterised by uncertainty. Considering the importance of many of these decisions, in particular at a strategic level, the amount of attention given to incorporating risk in the decision processes is surprisingly small. This may be partially explained by the complexity of optimisation models for these problems, as they include uncertain parameters, logical or other discrete decision variables, and more than one objective. Analytical tractability is hindered by this complexity, and even if mixed integer stochastic programming models can be developed, no efficient generic algorithms exist to solve them.

An important area where these issues are critical is the area of Capacity Expansion, in Operations Strategy. Capacity models are required to address a variety of problem features directly related to the abovementioned complexity. Partial or complete irreversibility of the investments, uncertainty in future rewards, some latitude on the timing or dynamics of the investments, multidimensionality of the invest-

ments, indivisibility of capacity expansions, fixed costs, economies of scale, the need to explicitly consider risk - the presence of these characteristics results, among other difficulties, in the presence of nonlinearities, nonconvexities, integrality and multiple objectives.

Metaheuristics are optimisation algorithms extremely well suited to efficiently tackle problems that present these features. In particular, multiobjective metaheuristics can be quite effective in simultaneously handling objectives reflecting both risk and, as traditionally, expected results. The application of multiobjective metaheuristics to mean-risk combinatorial optimisation problems is an unexplored research area, whose potential can thus be significant.

1.2 Research Strategy

Given the embryonic state of the research in this area, the primary objective of our work was to perform a preliminary assessment of multiobjective metaheuristics in solving mean-risk combinatorial optimisation problems.

This assessment was performed over a set of problems selected to represent the characteristics of the problems in the field:

- static single period structure, with all decisions made upfront, and dynamic multiperiod structure, with future opportunities for decisions that may use new information available up to the moment;
- binary, integer and mixed integer decision variables;
- linear, quadratic, closed-form and numerical objectives;
- expectation and classic (variance) and more recent (conditional value-at-risk) risk measures as objectives.

We have also tried to select problems exclusively related to Capacity Expansion, but were unsuccessful in finding a binary problem in this area. The three selected problems are the following:

- The *Static Stochastic Knapsack Problem*, a static binary problem, that we consider in an exact formulation, with numerical objectives, and in a sample approximation formulation, with closed-form linear and quadratic objectives. Both variance and conditional value-at-risk are considered as risk measures in both formulations.
- The *Multistage Capacity Investment Problem*, a dynamic integer problem, with uncertainty incorporated in the model by a scenario tree, and discrete capacity investment decision variables. Conditional value-at-risk is considered as a risk measure.
- The *Multistage Capacity Investment Problem with Process Flexibility*, a dynamic mixed integer problem, where uncertainty again is modelled by a scenario tree. Capacity investment decision variables are discrete, whereas capacity allocation decision variables are continuous. Conditional value-at-risk is again considered as a risk measure.

An object-oriented framework with multiobjective metaheuristics, that we have developed in previous work, was used in designing and implementing approaches for these problems. The ILOG CPLEX Mixed Integer Programming solver was used to obtain the mean-risk reference solution sets for randomly generated instances of all problems, except for the exact and variance formulations of the Static Stochastic Knapsack Problem, that required the use of a solution enumeration procedure. The performance of the developed metaheuristic approaches, in terms of solution quality and computational time, was then evaluated through a series of computational experiments.

1.3 Structure of the Dissertation

Chapter 2 is an introduction to some fundamental topics, aiming at providing the background for the research presented in this dissertation.

Chapters 3, 4 and 5 are related but self-contained essays that have been individually submitted for publication in international journals, currently being object of reviewing procedures. They have been slightly rearranged for inclusion in this dissertation, but the content has suffered no modifications. Chapter 3 is devoted to the Static Stochastic Knapsack Problem, chapter 4 to the Multistage Capacity Investment Problem, and chapter 5 to the Multistage Capacity Investment Problem with Process Flexibility.

Chapter 6 concludes the dissertation, presenting the key contributions of this research and suggesting future developments.

The option to provide self-contained essays in chapters 3, 4 and 5 entails some content repetition, namely of the following topics: description of the multiobjective local search template and overview of the object-oriented framework; description of performance measures; review of related work regarding optimisation with risk measures and metaheuristics for stochastic optimisation and portfolio selection; and definition of risk measures. This option, however, hopefully makes each of these chapters more accessible and readable.

Chapter 2

Some Fundamental Concepts, Tools and References

To help make this dissertation as self-contained as possible, we present in this chapter an introduction to some basic topics, and provide interested readers with a comprehensive set of relevant literature references. These topics are: fundamental concepts in multiobjective combinatorial optimisation; multiobjective metaheuristics, with an emphasis on multiobjective local search based approaches; the overall architecture of the MetHOOD object-oriented framework, highlighting the multiobjective local search template and the support for neighbourhood variation; an overview of the risk measures that are more commonly discussed in the literature and the prevailing concepts of adequacy of risk measures; a short review of results on capacity expansion, with some basic literature references, and additional references that suggest and validate the pertinence of the developments that we have proposed. A section is included for each of these topics.

2.1 Multiobjective Combinatorial Optimisation

A Combinatorial Optimisation (CO) problem is a mathematical optimisation problem with a discrete set of feasible solutions. It can be defined generically in the following way: given a discrete set S and a function $f : S \rightarrow \mathbb{R}$, find $\mathbf{x}^* \in S$ such that

$$f(\mathbf{x}^*) = \min \{f(\mathbf{x}) \mid \mathbf{x} \in S\}. \quad (2.1)$$

\mathbf{x} is a *feasible solution*, S is the *solution space* or *decision space* and f is the *objective function*. It is common and, in a certain way, natural to formulate a CO problem as an Integer Programming (IP) problem, in which the solutions are described by vectors of integer variables and the solution space is described by a set of equality and inequality constraints.

In terms of computational complexity, many CO problems are NP-hard, which reflects their intrinsic difficulty and justifies the adoption in practice of heuristic, non-optimising approaches.

Many practical problems, usefully modelled as CO problems, often require an evaluation of solutions according to a number of different perspectives. Multiobjective Combinatorial Optimisation (MOCO) problems can be represented by the following generic model:

$$\begin{aligned} \min \quad & f_1(\mathbf{x}) = z_1 \\ & \vdots \\ \min \quad & f_k(\mathbf{x}) = z_k \\ \text{s.t.} \quad & \mathbf{x} \in S, \end{aligned} \quad (2.2)$$

where \mathbf{x} is a *feasible solution*, S is the discrete *solution space*, and f_1, \dots, f_k are the *objective functions*. $\mathbf{z} = (z_1, \dots, z_k)$ is called a *criterion vector*. The feasible region in the objective space is $Z = \{\mathbf{z} \in \mathbb{R}^k : z_i = f_i(\mathbf{x}), \mathbf{x} \in S\}$ or $Z = \{\mathbf{z} \in \mathbb{R}^k : \mathbf{z} = \mathbf{f}(\mathbf{x}), \mathbf{x} \in S\}$, considering a vector function $\mathbf{f}(\mathbf{x}) = (f_1(\mathbf{x}), \dots, f_k(\mathbf{x}))$. $\mathbf{z} \in Z$ is *nondominated* if

and only if there is no other $\mathbf{z}' \in Z$ such that $z'_i \leq z_i, \forall i$, and $z'_i < z_i$, for some i . The *nondominated set* consists of all nondominated criterion vectors. $\mathbf{x} \in S$ is *efficient* if and only if its image in the objective space is nondominated. The *efficient set* consists of all efficient solutions.

It is often useful to work with similar ranges of values in all objectives. Such rescaling may be achieved by applying *range equalisation factors*. Accordingly, each objective z_i is multiplied by its corresponding range equalisation factor

$$\pi_i = \frac{1}{R_i} \left[\sum_{j=1}^k \frac{1}{R_j} \right]^{-1}, \quad (2.3)$$

where R_i is the *range width* of the i th criterion value over the efficient set.

As an important part of several methods for MOCO, *scalarising functions* can be used for mapping criterion vectors to values in an ordinal scale of quality. The *weighted sum* scalarising function $s_{ws}(\mathbf{z}, \mathbf{z}^0, \lambda) = \sum_{i=1}^k \lambda_i (z_i - z_i^0)$, considers a reference criterion vector \mathbf{z}^0 and strictly positive scalar weights λ_i . The *weighted sum problem* can be defined as

$$\begin{aligned} \min \quad & s_{ws}(\mathbf{z}, \mathbf{z}^0, \lambda) \\ \text{s.t.} \quad & \mathbf{z} \in Z. \end{aligned} \quad (2.4)$$

The optimal criterion vectors in this problem are designated as *supported nondominated*. Other nondominated criterion vectors are referred to as *nonsupported nondominated*. In linear multiobjective optimisation problems there are no nonsupported nondominated criterion vectors. However, in integer or nonlinear multiobjective problems, the existence of nonsupported nondominated criterion vectors is common.

For MOCO problems, the ϵ -*constraint* method considers single-objective problems constructed from the original multiobjective problem, where only one of the objective functions is kept as an objective, while the others are transformed into constraints

(that implicitly define their levels of achievement). By performing a systematic variation of the bounds of these constraints, this method is able to find both supported and nonsupported efficient solutions. For a bi-objective problem

$$\begin{aligned} \min \quad & f_1(\mathbf{x}) = z_1 \\ \min \quad & f_2(\mathbf{x}) = z_2 \\ \text{s.t.} \quad & \mathbf{x} \in S, \end{aligned} \tag{2.5}$$

we use the implementation of this method outlined in Algorithm 1, where δ is a positive constant small enough to avoid missing any efficient solutions.

Algorithm 1: ϵ -constraint method

```

Initialise the efficient set  $E = \{\}$ ;
 $\mathbf{x}_1 = \min f_1(\mathbf{x}), \text{ s.t. } \mathbf{x} \in S$ ;
while  $\mathbf{x}_1 \neq \{\}$  do
     $\mathbf{x}_2 = \min f_2(\mathbf{x}), \text{ s.t. } f_1(\mathbf{x}) = f_1(\mathbf{x}_1), \mathbf{x} \in S$ ;
    Insert  $\mathbf{x}_2$  in  $E$ ;
     $\mathbf{x}_1 = \min f_1(\mathbf{x}), \text{ s.t. } f_2(\mathbf{x}) \leq f_2(\mathbf{x}_2) - \delta, \mathbf{x} \in S$ ;
end

```

The use of *metrics*, for measuring distances between criterion vectors plays a fundamental role in several MOCO methods. The *Manhattan* metric defines the distance between two criterion vectors, \mathbf{z}^1 and \mathbf{z}^2 , by

$$\|\mathbf{z}^1 - \mathbf{z}^2\|_1^\pi = \sum_{i=1}^k \pi_i |z_i^1 - z_i^2|, \tag{2.6}$$

considering range equalisation factors π_i .

2.2 Multiobjective Metaheuristics

The difficulty in solving many practical CO problems has led to important research efforts in the development of more efficient approaches, reflected in significant reductions in computational requirements. Among these are *heuristics* - simple algorithms, frequently based in common sense, that are able to find a good (not necessarily optimal) solution for difficult problems, in a fast and easy way - and *metaheuristics* - master strategies that guide and modify other heuristics to produce solutions beyond those that can be produced by a search for local optima (Glover, 1986).

To a large extent, the success in applying metaheuristics to single-objective CO problems is due to features such as their general applicability, the flexibility to handle specific constraints in real world problems, and the interesting trade-off between solution quality and computation, development and implementation effort (Pirlot, 1996). Many of these methods also present a high robustness concerning the features of problem instances or parameter tuning. Tabu Search (TS), Simulated Annealing (SA) and Genetic Algorithms (GA) are metaheuristics that are currently broadly used, and described in mainstream Operations Research textbooks.

These same features have been fostering their application in MOCO, enabling the handling of variations in problem formulation or in the objectives. Surveys on multiobjective metaheuristics (MOMH) are available in Ehrgott and Gandibleux (2000), Jones et al. (2002) and Ehrgott and Gandibleux (2004). GA have led the way in this area, the pioneer work of Schaffer with the Vector Evaluated Genetic Algorithm (Schaffer, 1985) dating back to 1985. The early survey on multiobjective GA in Fonseca and Fleming (1995) points out that GA seem specially fit for use in multiobjective contexts for two main reasons: working with a population of solutions, they can search for the multiple solutions of the efficient set in parallel, eventually exploring similarities among them; also, they are less sensitive than traditional mathematical

programming techniques to the issues of shape and continuity of the nondominated set.

This second feature is shared with SA and TS based approaches. One group of these approaches (Serafini, 1992; Ulungu et al., 1998; Gandibleux et al., 1997) is based on repeated executions of the single-objective metaheuristic, with a combination of the objectives in an aggregating function, usually a weighted sum scalarising function. A search direction is established by the weights in this function, whose variation in each execution aims at enabling a complete approximation of the nondominated set. In another group, that includes Pareto Simulated Annealing (PSA) (Czyzak and Jaskiewicz, 1998) and Tabu Search for Multiobjective Combinatorial Optimisation (TAMOCO) (Hansen, 2000), the first mentioned feature is introduced in SA and TS based approaches, through the consideration of a population of solutions, each one holding its own set of weights. These weights are dynamically computed so that each solution moves towards the nondominated frontier and away from other solutions in the population, that are nondominated with respect to it.

In PSA (Algorithm 2) the weights for each solution are updated according to the relation between the components of the corresponding criterion vector and the nearest nondominated criterion vector. The distance between criterion vectors may be measured with a Manhattan metric. A constant multiplying factor α , or its inverse $1/\alpha$, are used to update the weights, with α higher than, but close to, 1 (e.g., 1.05). The weights are incorporated in the acceptance probability, that can be defined as the minimum of the weighted acceptance probabilities for each objective

$$\min_{j=1,\dots,k} \left\{ \min \{1, \exp(\Delta z_{i,j}/T)\}^{\lambda_{i,j}} \right\}. \quad (2.7)$$

Increasing the weight associated to an objective reduces the probability of accepting movements that do not improve that objective and increases the probability of im-

Algorithm 2: Pareto Simulated Annealing

```

Generate a set of initial feasible solutions  $G \subset S$ ;
Initialise the approximation to the efficient set  $E = \{\}$ ;
foreach  $\mathbf{x}_i \in G$  do
    Update  $E$  with  $\mathbf{x}_i$ ;
end
Initialise temperature  $T$ ;
while a stopping criterion is not met do
    foreach  $\mathbf{x}_i \in G$  do
        Select solution  $\mathbf{x}_w \in G$ , such that  $\mathbf{f}(\mathbf{x}_w)$  is nearest to and nondominated by  $\mathbf{f}(\mathbf{x}_i)$ ;
        if  $\mathbf{x}_w$  does not exist or first iteration with  $\mathbf{x}_i$  then
            Generate random weights  $\lambda_{i,j}$  for the weight vector  $\lambda_i$  associated with  $\mathbf{x}_i$ , such
            that  $\lambda_{i,j} \geq 0$  and  $\sum_{j=1}^k \lambda_{i,j} = 1$ ;
        end
        else
            foreach objective function  $f_j$  do
                if  $f_j(\mathbf{x}_i) \leq f_j(\mathbf{x}_w)$  then  $\lambda_{i,j} = \alpha \lambda_{i,j}$ ;
                else  $\lambda_{i,j} = \lambda_{i,j} / \alpha$ ;
            end
            Normalise weights  $\lambda_{i,j}$  so that  $\sum_{j=1}^k \lambda_{i,j} = 1$ ;
        end
        Randomly select  $\mathbf{x}_s \in \text{Neighbourhood}(\mathbf{x}_i)$ ;
        if  $\mathbf{f}(\mathbf{x}_s)$  is nondominated by  $\mathbf{f}(\mathbf{x}_i)$  then update  $E$  with  $\mathbf{x}_s$ ;
        Randomly select a value  $p \in [0; 1]$ ;
        if  $p \leq P(\mathbf{f}(\mathbf{x}_i), \mathbf{f}(\mathbf{x}_s), T, \lambda_i)$  then  $\mathbf{x}_i = \mathbf{x}_s$ ;
    end
    Update temperature  $T$ ;
end

```

proving that objective. The temperature starts at an initial value T_0 and, at every L iterations, is multiplied by a constant positive value lower than 1.

In TAMOCO (Algorithm 3) the vector of weights is used to define a direction of optimisation for each solution, towards the nondominated set and away from other nondominated solutions, in proportion to their *proximity*. Proximity can be defined as the inverse of distance, which in turn can be measured with a Manhattan metric, considering range equalisation factors. In the absence of better knowledge, range equalisation factors can be computed from the objective ranges in the approximation set.

Algorithm 3: Tabu Search for Multiobjective Combinatorial Optimisation

```

Generate a set of initial feasible solutions  $G \subset S$ ;
Initialise the vector of range equalisation factors  $\pi$  with  $\pi_j = 1/k$ ;
Initialise the approximation to the efficient set  $E = \{\}$ ;
foreach  $\mathbf{x}_i \in G$  do
    Initialise the corresponding tabu list  $TL_i = \{\}$ ;
    Update  $E$  with  $\mathbf{x}_i$ ;
end
while a stopping criterion is not met do
    foreach  $\mathbf{x}_i \in G$  do
        Initialise the corresponding weight vector  $\lambda_i = \mathbf{0}$ ;
        foreach  $\mathbf{x}_1 \in G$  such that  $\mathbf{f}(\mathbf{x}_1)$  is nondominated by and different from  $\mathbf{f}(\mathbf{x}_i)$  do
            Compute proximity  $w = g(d(\mathbf{f}(\mathbf{x}_1), \mathbf{f}(\mathbf{x}_i)), \pi)$ ;
            foreach objective function  $f_j$  such that  $f_j(\mathbf{x}_i) < f_j(\mathbf{x}_1)$  do
                 $\lambda_{i,j} = \lambda_{i,j} + \pi_j w$ ;
            end
            Normalise weights  $\lambda_{i,j}$  so that  $\sum_{j=1}^k \lambda_{i,j} = 1$ ;
        end
        if  $\lambda_i = \mathbf{0}$  then
            Generate random weights  $\lambda_{i,j}$  for the weight vector  $\lambda_i$  associated with  $\mathbf{x}_i$ , such
            that  $\lambda_{i,j} \geq 0$  and  $\sum_{j=1}^k \lambda_{i,j} = 1$ ;
        end
        Select  $\mathbf{x}_s \in \text{Neighbourhood}(\mathbf{x}_i)$ , such that
         $\lambda_i \cdot \mathbf{f}(\mathbf{x}_s) \leq \lambda_i \cdot \mathbf{f}(\mathbf{x}_t), \forall \mathbf{x}_t \in \text{Neighbourhood}(\mathbf{x}_i)$ , and
        ( $TL_i$  does not make  $(\mathbf{x}_i, \mathbf{x}_s)$  tabu or  $\mathbf{x}_s$  satisfies the aspiration criterion);
        Insert attributes of  $(\mathbf{x}_i, \mathbf{x}_s)$  at the end of  $TL_i$ , removing the first element if  $TL_i$  is
        full;
         $\mathbf{x}_i = \mathbf{x}_s$ ;
        Update  $E$  with  $\mathbf{x}_s$ ;
        Update  $\pi$ ;
    end
    if a drift criterion is met then
        Replace a randomly selected solution with another randomly selected solution in  $G$ ;
    end
end

```

To avoid the concentration of solutions in certain areas, a *drift* mechanism is used, whereby a randomly selected solution is replaced with a copy of another randomly selected solution. The aspiration criterion consists of accepting any efficient solution.

PSA and TAMOCO can be viewed as Multiobjective Local Search (MOLS) procedures. Both aim at producing a good approximation of the efficient set, working with a population of solutions, each solution holding a weight vector for the definition

Algorithm 4: Multiobjective Local Search Template

```

Generate a set of initial feasible solutions  $G \subset S$ ;
Initialise the approximation to the efficient set  $E = \{\}$ ;
foreach  $\mathbf{x}_i \in G$  do
    Initialise the corresponding context;
    Update  $E$  with  $\mathbf{x}_i$ ;
end
while a stopping criterion is not met do
    foreach  $\mathbf{x}_i \in G$  do
        Update the corresponding weight vector  $\lambda_i$ ;
        Initialise the selected solution  $\mathbf{x}_s = 0$ ;
        foreach  $\mathbf{x}' \in \text{Neighbourhood}(\mathbf{x}_i)$  do
            Update  $E$  with  $\mathbf{x}'$ ;
            if  $\mathbf{x}'$  is selectable and  $\mathbf{x}'$  is preferable to  $\mathbf{x}_s$  then  $\mathbf{x}_s = \mathbf{x}'$ ;
        end
        if  $\mathbf{x}_s \neq 0$  and  $\mathbf{x}_s$  is acceptable then  $\mathbf{x}_i = \mathbf{x}_s$ ;
    end
end
  
```

of a search direction. Each approach proposes a different strategy for the definition of the weights, but share identical purposes for that definition: orientation of the search towards the nondominated frontier and spreading of solutions over that frontier (the former is achieved by the use of positive weights, while the latter is based on a comparison with other solutions of the population). Although in different ways, both methods operate on each single solution, searching and selecting a solution in its neighbourhood that will eventually replace it. Each procedure involves traditional metaheuristic components such as neighbourhoods, in general, or tabu lists, in the specific case of TAMOCO. The identification of these common aspects has suggested the definition of a MOLS template (Algorithm 4), as a way to provide a generic basis for designing different specific algorithms.

2.3 MetHOOD - a Metaheuristic Framework

2.3.1 Object-Oriented Software

Four essential principles of the object-oriented paradigm are *data encapsulation*, *data abstraction*, *inheritance* and *polymorphism*. The *object* is the unit of data encapsulation, consisting of a set of *variables* and a set of *methods* used to alter and access those variables. An object accepts *messages* that invoke methods. An object's *signature* is the set of messages it accepts. A *class* is a data abstraction, expressing aspects that are common to identical objects, and taking the form of a template to create a particular kind of object. Class inheritance allows a set of classes to share parts of a common signature or implementation (variables and methods). Polymorphism is essential to code reuse, by enabling methods to take objects of different types as arguments.

Reusable object-oriented software has been made available mainly through *class libraries* and *frameworks*. A class library packages together a set of classes, eventually structured using inheritance, from which an application can be built. A framework can be defined as a set of classes that embodies an abstract design for solutions to a family of related problems (Johnson and Foote, 1988). One main difference between these two concepts is that frameworks provide default behaviour, while with class libraries, all the collaboration between components usually has to be defined. Frameworks provide design reuse, an area where another domain has also gained particular significance - *Design Patterns*, which are descriptions of communicating objects and classes that are customised to solve a general design problem in a particular context (Gamma et al., 1994).

2.3.2 Object-Oriented Approaches for Metaheuristics

The two main incentives for developing object-oriented approaches for metaheuristics have been bringing theory and applications closer, by developing simply structured, open-ended systems, that incorporate theoretical results (Nievergelt, 1994), and facilitating the implementation and comparison of methods, through increased modularity, rational order and reusability of software structures (Ferland et al., 1996). Several object-oriented approaches for metaheuristics have been presented in the literature. Descriptions of some of the most prominent can be found, in detail, in Voss and Woodruff (2002) or, in a brief introduction, in Fink et al. (2002).

2.3.3 MetHOOD

MetHOOD (**M**eta**H**euristics **O**bject-**O**riented **D**evelopment) is a framework for MOMH that extensively incorporates design patterns in its design (Claro and Sousa, 2001). At the time the MetHOOD framework was proposed, no other object-oriented approaches for MOMH had been reported in the literature. Still in 2001, a C++ class library for MOMH, developed by Andrzej Jaszkievicz, was made publicly available at <http://www-idss.cs.put.poznan.pl/~jaszkiewicz/MOMHLib/>. In a 2004 review of heuristics and metaheuristics designed for the solution of MOCO problems (Ehrgott and Gandibleux, 2004) MetHOOD was still the only work cited in the area of reusable MOMH software.

MetHOOD has been used to support the applications described in the following chapters and in its present state of development, provides (Figure 2.1): support for the definition of the variable parts of the problem domain, related to solutions, movements, increments and evaluating functions; support for problem data; a template and a concrete implementation of a constructive algorithm; a MOLS template and the derivations of PSA and TAMOCO from this template; extensions for a candidate

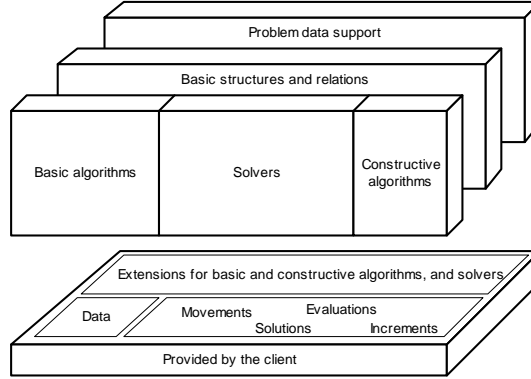


Figure 2.1: The MetHOOD framework

list strategy and a neighbourhood variation strategy; a solver level for the articulation of constructive and MOLS algorithms, and the implementation of a high-level, parallel, hybridisation strategy.

2.3.4 MOLS Template

Figure 2.2 presents the class diagram for the MOLS template part of the framework. A MOLS algorithm (*MOLocalSearch*) iterates over a population ($Population_{Initial}$) of solutions (*Solution*), building an approximation to the efficient set (*EfficienSet*) of the considered problem. A neighbourhood (*Neighbourhood*) is created for each solution, with the services of a neighbourhood prototype ($Neighbourhood_{Prototype}$). For each solution, the algorithm obtains all the movements ($Movement_{Current}$) in the neighbourhood and selects one ($Movement_{Selected}$). The neighbourhood is traversed with a movement iterator (*MoveIterator*), that sees the neighbourhood as a move-container (*MoveContainer*). The efficient set is updated with all generated movements. Movement selection always involves the solution's weights, which are defined by a weight definition strategy (*WeightDefinition*), taking into consideration the solutions from a larger population ($Population_{Larger}$). This larger population

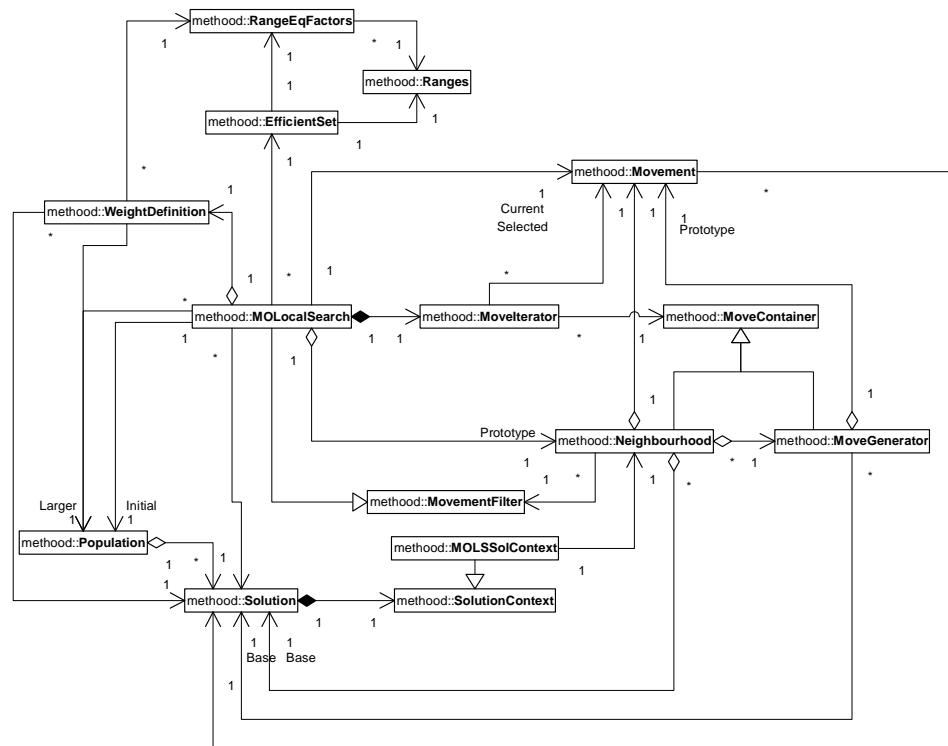


Figure 2.2: MOLS class diagram

contains the population of the algorithm, and may also contain other solutions, such as reference solutions, or solutions being worked on by other algorithms.

2.3.5 Implementation of PSA and TAMOCO

Figure 2.3 illustrates the differences in the way that PSA (*MOSimAnneal*) and TAMOCO (*MOTabuSearch*) implement the definition of several of the template’s primitive operations: weight vectors are distinctly initialised and updated (for PSA *PSAWeightDef*, and for TAMOCO *MOTSWeightDef*); the neighbourhood in PSA is a random subneighbourhood with just one movement; the selection of a generated movement (*+IsMovementValid()*) in TAMOCO considers tabu status, aspiration criteria, and a comparison of evaluations based on a weighted sum scalarising

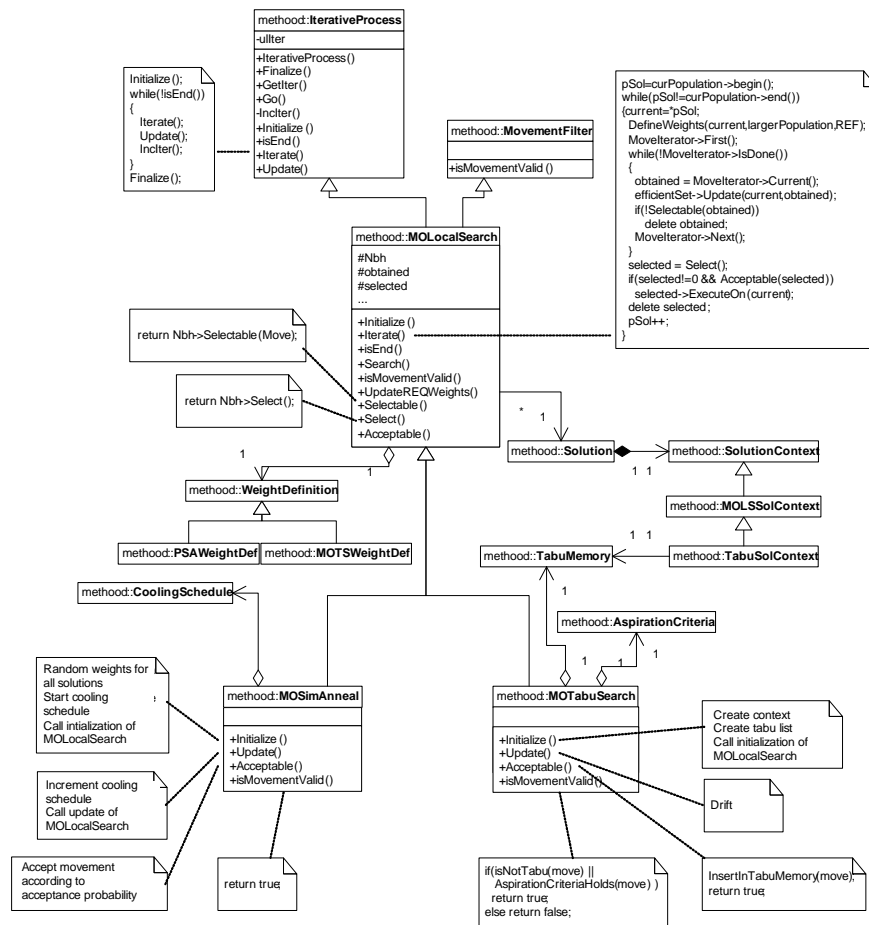


Figure 2.3: PSA and TAMOCO derived from the MOLS template

function, while in PSA a generated movement is always selected; the acceptance ($+Acceptable()$) of a selected movement in PSA is a function of an acceptance probability, while in TAMOCO a selected movement is always accepted.

2.3.6 Neighbourhood Variation

The framework also provides support for neighbourhood variation (Figure 2.4), by considering a sequence of neighbourhood structures (*vector* $< MoveGenerator >$) and using them dynamically according to the evolution of the search process: if a new

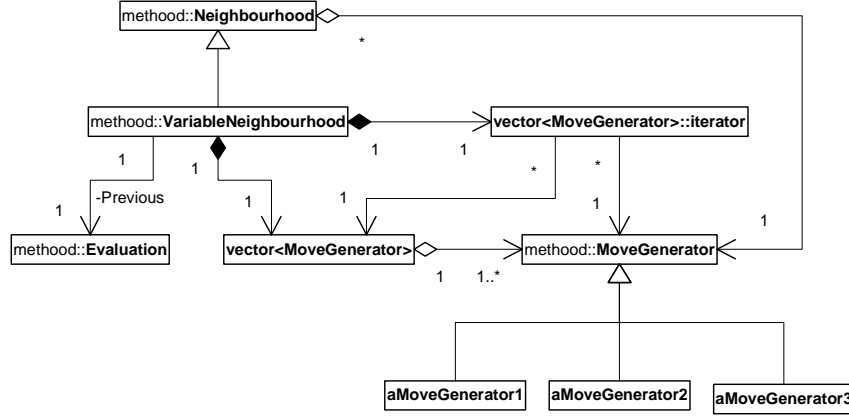


Figure 2.4: Class diagram for neighbourhood variation

accepted solution is preferable to the current one, or if the current neighbourhood is the last in the sequence, the first neighbourhood in the sequence will be used next; otherwise the following neighbourhood in the sequence will be used next.

2.4 Risk Measures

2.4.1 Definitions of Risk Measures

The identification of adequate risk measures is currently a field of very active research, that we will not review in detail in this work. This section is a very short introduction to the basic concepts of risk measurement, aiming at providing an adequate context for the work that is presented ahead.

Measuring risk requires that a correspondence ρ is established between a space \mathbb{X} of random variables and a nonnegative real number, i.e., $\rho : \mathbb{X} \rightarrow \mathbb{R}_0^+$. Scalar measures of risk allow ordering and comparison according to risk values. Most of the basic ideas for risk measures arise from the consideration of dispersion parameters, excess probabilities, quantiles or conditional expectations.

For a presentation of some of the fundamental concepts in risk measures, we consider a continuous loss random variable X with distribution function F_X and density function f_X . The *expected value* of X can be defined in the following way

$$E[X] = \int x f_X(x) dx. \quad (2.8)$$

Table 2.1 presents some of the risk measures that are more commonly discussed in the literature $((\cdot)_+ = \max\{\cdot, 0\})$.

Table 2.1: Definitions of risk measures

Risk measure	Definition
Variance	$E[(X - E[X])^2]$
Mean absolute deviation	$E[X - E[X]]$
p -th semideviation from target t	$(E[(X - t)_+^p])^{1/p}$
p -th central semideviation	$(E[(X - E[X])_+^p])^{1/p}$
Gini mean difference	$\int E[(\xi - X)_+] dF_X(\xi)$
α -value-at-risk (VaR_α)	$\inf\{x F_X(x) > \alpha\}$
α -conditional value-at-risk (CVaR_α)	$E[X X \geq \text{VaR}_\alpha[X]]$

To be able to handle possible discontinuities, the definition of CVaR_α must replace the original conditional expectation by the following α -tail expectation

$$\begin{aligned} \text{CVaR}_\alpha[X] &= \int x dF_X^\alpha(x), \\ \text{where } F_X^\alpha(x) &= \begin{cases} 0 & \text{if } F_X(x) < \alpha \\ \frac{F_X(x) - \alpha}{1 - \alpha} & \text{if } F_X(x) \geq \alpha \end{cases}. \end{aligned} \quad (2.9)$$

The prevailing concepts of adequacy of risk measures are *consistency with stochastic dominance* (Ogryczak and Ruszczyński, 1999) and *coherence* (Artzner et al., 1999).

2.4.2 Stochastic Dominance

The stochastic dominance (Ogryczak and Ruszczyński, 1999) relations between two random variables are defined by pointwise comparison of performance functions based on their distribution functions. The first function $F_X^{(1)}$ is just the distribution function $F_X^{(1)}(x) = F_X(x)$, and the weak (\succeq) and strict (\succ) relations of the first degree stochastic dominance (FSD) are defined as follows:

$$\begin{aligned} X \succeq_{\text{FSD}} Y &\Leftrightarrow F_X(z) \geq F_Y(z), \forall z \in \mathbb{R}. \\ X \succ_{\text{FSD}} Y &\Leftrightarrow X \succeq_{\text{FSD}} Y \text{ and } Y \not\succeq_{\text{FSD}} X. \end{aligned} \quad (2.10)$$

$F_X(x)$ expresses the probability of having losses below a given target value x . If $X \succ_{\text{FSD}} Y$, then X is preferred to Y within all models preferring smaller losses, regardless of risk-aversion.

The second function $F_X^{(2)}$ is given by the area below the distribution function $F_X(x)$:

$$F_X^{(2)}(x) = \int_{-\infty}^x F_X(z) dz, x \in \mathbb{R}. \quad (2.11)$$

The weak and strict relations of the second degree stochastic dominance (SSD) are defined as follows:

$$\begin{aligned} X \succeq_{\text{SSD}} Y &\Leftrightarrow F_X^{(2)}(z) \geq F_Y^{(2)}(z), \forall z \in \mathbb{R}. \\ X \succ_{\text{SSD}} Y &\Leftrightarrow X \succeq_{\text{SSD}} Y \text{ and } Y \not\succeq_{\text{SSD}} X. \end{aligned} \quad (2.12)$$

If $X \succ_{\text{SSD}} Y$, then X is preferred to Y within all risk-averse preference models that prefer smaller losses. Consistency with SSD is therefore a fundamental aspect in risk comparison.

Consistency with stochastic dominance has been studied in the literature mainly considering the following definition: a risk measure ρ is α -consistent with stochastic

dominance of order p iff

$$X \succeq_p Y \Rightarrow E(X) + \alpha \rho(X) \leq E(Y) + \alpha \rho(Y). \quad (2.13)$$

The α -consistency with stochastic dominance of order p implies α' -consistency with stochastic dominance of order p , for all α' such that $0 < \alpha' \leq \alpha$ (Ogryczak and Ruszczyński, 1999). A risk measure ρ is consistent with stochastic dominance of order p iff it is α -consistent with stochastic dominance of order p for all $\alpha \in \mathbb{R}^+$.

In general, variance and mean absolute deviation are not consistent with FSD or SSD (see, e.g., Märkert (2004)). p -th central semideviation is 1-consistent with stochastic dominance of order $1, \dots, p+1$ (Ogryczak and Ruszczyński, 1999) and the same result applies to Gini's mean difference (Yitzhaki, 1982). p -th semideviation from target t is consistent with FSD and SSD, for all $p \in \mathbb{N}$ (Fishburn, 1977). VaR is consistent with FSD, but not SSD, whereas CVaR is consistent with both (Ogryczak and Ruszczyński, 2002).

2.4.3 Coherence

A risk measure $\rho : \mathbb{X} \rightarrow \mathbb{R}$ is coherent (Artzner et al., 1999) if it satisfies the following properties:

- *Translation invariance*: $\rho(X + a) = \rho(X) + a, \forall X \in \mathbb{X}, \forall a \in \mathbb{R}$.
- *Subadditivity*: $\rho(X + Y) \leq \rho(X) + \rho(Y), \forall X, Y \in \mathbb{X}$.
- *Monotonicity*: if $X \leq Y$, then $\rho(X) \geq \rho(Y), \forall X, Y \in \mathbb{X}$.
- *Positive homogeneity*: $\rho(\lambda X) = \lambda \rho(X), \forall X \in \mathbb{X}, \forall \lambda > 0$.

In general, variance, mean absolute deviation, p -th central semideviation and Gini's mean difference are not coherent, and p -th semideviation from target t requires

adaptation of the target to the random variable to be coherent (Märkert, 2004). VaR is not coherent (Artzner et al., 1999). CVaR is coherent (Rockafellar and Uryasev, 2002).

2.4.4 Computing CVaR_α via Linear Programming

The properties referred above partially justify the increasing attention that CVaR_α has been receiving in the literature. Another important reason for this attention is the fact that it can be computed via linear programming. The problem of minimising CVaR_α can be formulated as follows (Rockafellar and Uryasev, 2002):

$$\begin{aligned} \min \quad & \xi + \frac{1}{(1-\alpha)N} \sum_{j=1}^N Z_j \\ \text{s.t.} \quad & Z_j \geq f(\mathbf{x}, \omega_j) - \xi, \quad j = 1, \dots, N \\ & \mathbf{x} \in S, Z_j \geq 0, \xi \geq 0, \end{aligned} \tag{2.14}$$

where \mathbf{x} is a solution, S is the solution space, ω is the randomness component with a certain probability distribution, ω_j are scenarios of ω , $j = 1, \dots, N$, and $f(\mathbf{x}, \omega)$ is a loss function. If $f(\mathbf{x}, \omega)$ is linear and S is described with linear constraints, the problem of minimising CVaR_α is a linear programming problem.

CVaR_α constraints can also be considered, as in the following formulation for the minimisation of the mean loss, with a bound of C for the CVaR_α :

$$\begin{aligned} \min \quad & \frac{1}{N} \sum_{j=1}^N f(\mathbf{x}, \omega_j) \\ \text{s.t.} \quad & \xi + \frac{1}{(1-\alpha)N} \sum_{j=1}^N Z_j \leq C \\ & Z_j \geq f(\mathbf{x}, \omega_j) - \xi, \quad j = 1, \dots, N \\ & \mathbf{x} \in S, Z_j \geq 0, \xi \geq 0. \end{aligned} \tag{2.15}$$

2.5 Capacity Expansion

2.5.1 Background

A recent review of capacity expansion literature (Julka et al., 2007) identifies the first recognition of the importance of capacity expansion in Operations Strategy in Wheelwright (1978), where it was regarded as one of the five strategic manufacturing decision areas, and cites Rudberg and Olhager (2003) to report that substantial subsequent research has provided wide support for this view. This decision area still remains crucial, particularly for manufacturing corporations with global production facilities (Julka et al., 2007) and in high-tech industries such as semiconductor, consumer electronics, telecommunications and pharmaceutical (Wu et al., 2005).

In a 2003 landmark review on strategic capacity management under uncertainty (Van Mieghem, 2003), capacity expansion has been defined as being concerned with deciding the type, magnitude, timing, and location of capacity acquisition: these decisions about processing resources in a network play a fundamental role in defining the network's capabilities, and are associated with decisions on the types and the levels of investment. Three very important characteristics are present in investment, in varying degrees (Dixit and Pindyck, 1994): partial or complete irreversibility, uncertainty over the future rewards and some latitude on the timing or dynamics of the investment. A fourth characteristic is added in Van Mieghem (2003): multidimensionality, i.e., the possibility of investing in resources with different financial and operational properties. This review also points out additional fundamental challenges in capacity cost modelling - indivisibility of capacity expansions and nonconvexity arising, e.g., from fixed costs or economies of scale - and the surprising fact that few papers on capacity investment tackle issues related to risk, even if we are often facing significant investments with uncertain future rewards.

2.5.2 Relevant Literature

Recently the research on capacity expansion has been concerned with the development of models and techniques that are able to deal with this difficult set of characteristics. We have found in Ahmed et al. (2003) and Ahmed and Sahinidis (2003) fundamental references for our work, since the authors have put forward capacity expansion models that incorporate several of the characteristics pointed out above: irreversibility, uncertainty, latitude on the timing of investments, multidimensionality and nonconvex cost functions.

Those authors have divided the previous relevant literature in three main groups:

- *Early approaches based on stochastic control theory.* These approaches use simple stochastic processes to model demand, for analytical tractability. Manne (1961) is the first reference on dynamic capacity models with stochastic demand. Other references are Freidenfelds (1980), Davis et al. (1987) and Bean et al. (1992).
- *Two-stage stochastic programming.* Eppen et al. (1989) use standard mixed integer programming to solve a two-stage stochastic programming model with fixed charge expansion cost functions, incorporating elements of scenario planning, integer programming and risk analysis, for strategic capacity planning in the automotive industry. Berman et al. (1994) apply a two stage stochastic model with linear costs to capacity expansions in services, using Lagrangian relaxation. Fine and Freund (1990) formulate and study a product-flexible capacity investment model as a two-stage nonlinear stochastic program, but assuming linearity in the cost functions. Liu and Sahinidis (1996) propose a two-stage stochastic programming approach for process planning under uncertainty, extending a deterministic mixed-integer linear programming formulation to account for the presence of discrete random parameters and subsequently devising a decompo-

sition algorithm for the solution of the stochastic model. Swaminathan (2000) provides heuristics for a two-stage model applied to tool capacity planning in the semiconductor industry, under uncertainty in demand and with capacity decisions in the first stage.

- *Multistage stochastic programming.* Rajagopalan et al. (1998) develop a dynamic programming algorithm for a multi-stage capacity acquisition and replacement problem, where capacity availability is anticipated, but its magnitude and timing are uncertain. Chen et al. (2002) use Lagrangian decomposition to solve a problem of multistage stochastic capacity expansion with technology selection.

2.5.3 A Multistage Stochastic Integer Programming Model

Addressing the limitations identified in this body of research, those authors propose a multistage stochastic integer programming model, with a scenario tree to model the stochastic evolution of costs and demand, and fixed-charge cost functions to model economies of scale. This model is tackled with approximation and reformulation schemes that lead to significant improvements on the computational times obtained by a straightforward use of IP solvers.

The more generic model presented in Ahmed and Sahinidis (2003), addresses the problem of determining the timing and the level of capacity acquisitions for a set of production facilities \mathcal{I} , as well as an allocation of capacity to satisfy the demand of a set of product families \mathcal{J} . The capacity expansion and allocation decisions are made with the objective of minimising the expected total discounted investment and allocation cost, for a discretised planning horizon.

The product demands (d), fixed and variable costs of capacity acquisition (α and β), and the costs of allocating capacity to products (γ) are assumed to be stochastic.

Uncertainty is modelled as a multilayered tree, whose levels correspond to time periods. The nodes at a certain tree level constitute the states of the world that can be distinguished by information available up to that period. $\mathcal{T}(n)$ denotes the subtree rooted in node n , with $n = 0$ being the root node, and $\mathcal{P}(n)$ the path from the root node to node n . The probability associated with the state of the world in node n is p_n .

\mathcal{S} is the set of leaf nodes, each related to one of the S scenarios. A scenario corresponds to a path from the root to a leaf, representing a joint realisation of the uncertain parameters over all time periods, i.e., for the scenario corresponding to a leaf node $m \in \mathcal{S}$, $\{d_{j,n}, \alpha_{i,n}, \beta_{i,n}, \gamma_{i,j,n}\}_{n \in \mathcal{P}(m)}$, where $i \in \mathcal{I}$ and $j \in \mathcal{J}$.

Figure 2.5 presents an example of a binary scenario tree, displaying the expansion costs for a problem with 4 facilities and 3 time periods. In each node the left column shows the fixed costs, and the right column the variable costs. Each row corresponds to a different facility. To each arc is associated the probability of the node where the arc is directed to.

In each node $n \in \mathcal{T}(0)$, each facility $i \in \mathcal{I}$ is characterised by a discounted fixed cost $\alpha_{i,n}$ and a discounted variable cost $\beta_{i,n}$ of expansion. The capacity expansions are (deterministically) bounded by $M_{i,n}$. The initial capacities are zero, but the adaptations to include initial capacities are straightforward. The demand for product $j \in \mathcal{J}$ is given, in each node, by $d_{j,n}$. Each unit of capacity of facility i can produce $q_{i,j}$ units of product j , and the discounted cost associated with this allocation in node n is $\gamma_{i,j,n}$.

The decision variables are $x_{i,n}$, the capacity expansion for facility i at node n , and $w_{i,j,n}$, the number of units of capacity of facility i allocated to the production of product j in node n . The binary variables $y_{i,n}$ take the value 1 if the capacity of facility i is expanded in node n , and the value 0 otherwise.

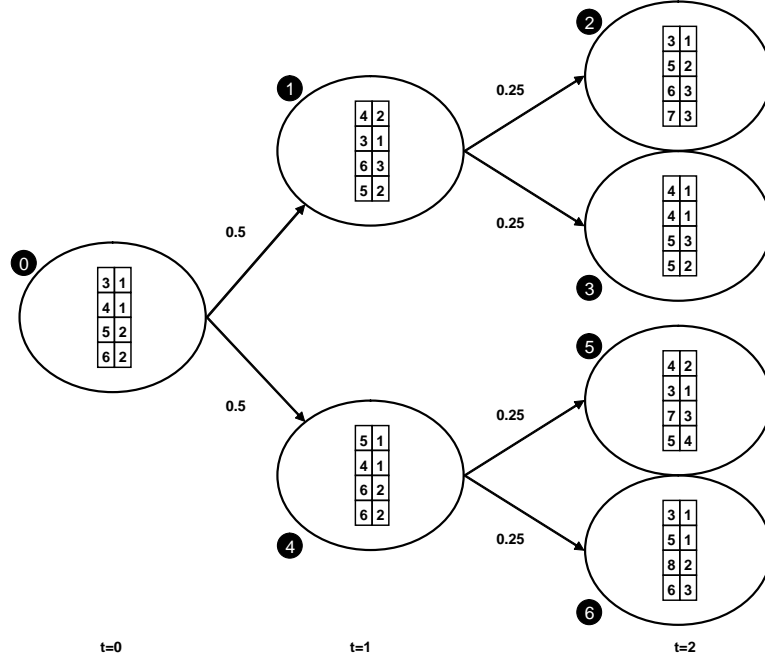


Figure 2.5: Scenario tree for expansion costs

This problem can be formulated as follows:

$$\begin{aligned}
 \min \quad & \sum_{m \in \mathcal{S}} \left(p_m \sum_{n \in \mathcal{P}(m)} \sum_{i \in \mathcal{I}} \left(\alpha_{i,n} y_{i,n} + \beta_{i,n} x_{i,n} + \sum_{j \in \mathcal{J}} (\gamma_{i,j,n} w_{i,j,n}) \right) \right) \\
 \text{s.t.} \quad & x_{i,n} \leq M_{i,n} y_{i,n}, \quad n \in \mathcal{T}(0), i \in \mathcal{I} \\
 & \sum_{i \in \mathcal{I}} q_{i,j} w_{i,j,n} \geq d_{j,n}, \quad n \in \mathcal{T}(0), j \in \mathcal{J} \\
 & \sum_{j \in \mathcal{J}} w_{i,j,n} \leq \sum_{o \in \mathcal{P}(n)} x_{i,o}, \quad n \in \mathcal{T}(0), i \in \mathcal{I} \\
 & x_{i,n} \geq 0, \quad n \in \mathcal{T}(0), i \in \mathcal{I} \\
 & y_{i,n} \in \{0, 1\}, \quad n \in \mathcal{T}(0), i \in \mathcal{I} \\
 & w_{i,j,n} \geq 0, \quad n \in \mathcal{T}(0), i \in \mathcal{I}, j \in \mathcal{J}.
 \end{aligned} \tag{2.16}$$

The first set of constraints define $y_{i,n}$ in terms of variables $x_{i,n}$ and establish the

bounds for capacity expansions. The second and third sets of constraints are the demand satisfaction and capacity constraints, respectively.

2.5.4 Challenges in Capacity Expansion

Lumpiness and the explicit consideration of risk are critical practical issues mentioned in Van Mieghem (2003) that are not properly addressed in the literature, and whose importance is reinforced by the survey on capacity management in high-tech industries by Wu et al. (2005), that confirms the relevance of those features for capacity decisions in such industries.

In the most recent review of capacity expansion literature, Julka et al. (2007) identify a primary research opportunity in developing models to simultaneously handle the multiple factors that are relevant in the decision making processes involved in capacity expansion. As Van Mieghem (2003) concludes, "Fortunately and unfortunately, capacity-portfolio models rapidly become complex. Complexity is unfortunate because it often makes superior analytical solutions elusive. Thus, simulation-based optimization becomes the natural second-best option and is expected to increase in popularity. At the same time, complexity is fortunate as study is worthwhile with a potential impact on practice. Compared to the impact of financial portfolio analysis, even a fraction would be substantial."

Chapter 3

A Multiobjective Metaheuristic for a Mean-Risk Static Stochastic Knapsack Problem

(Under review at Computational Optimization and Applications)

In this paper we address two major challenges presented by stochastic discrete optimisation problems: the multiobjective nature of the problems, once risk aversion is incorporated, and the frequent difficulties in computing exactly, or even approximately, the objective function. The latter has often been handled with methods involving sample average approximation, where a random sample is generated so that population parameters may be estimated from sample statistics - usually the expected value is estimated from the sample average. We propose the use of multiobjective metaheuristics to deal with these difficulties, and apply a multiobjective local search metaheuristic to both exact and sample approximation versions of a mean-risk static stochastic knapsack problem. Variance and conditional value-at-risk are considered as risk measures. Results of a computational study are presented, that indicate the

approach is capable of producing high-quality approximations to the efficient sets, with a modest computational effort.

3.1 Introduction

A large number of decisions in Operations Management are made in the presence of uncertainty. In fact, key factors, such as prices, resource availability or product demand, are regularly characterised by uncertainty. Considering the importance of many of these decisions, in particular at a strategic level, the amount of attention given to incorporating risk in the decision processes is surprisingly small. This may be partially explained by the complexity of optimisation models for these problems, as they include uncertain parameters, logical or other discrete decision variables, and more than one objective.

Even if these problems can be formulated as mixed integer stochastic programming problems, no efficient generic algorithms exist to solve them, in spite of the recent increase in the attention given to integrality in the stochastic programming literature. Research on the application of metaheuristics to these problems, on the other hand, has either focused on single objective problems or had very confined applications, particularly in the areas of robust optimisation and portfolio selection.

In this paper we perform a preliminary assessment of multiobjective metaheuristics for tackling stochastic combinatorial optimisation problems, by applying a multiobjective local search metaheuristic to a problem with the previously mentioned difficulties - the static stochastic knapsack problem - that we cast in a mean-risk framework. We use two different risk measures - variance and conditional value-at-risk - and consider an exact version of the problem, where expectation and risk measures are computed exactly, and a sample approximation version, where those values are computed from a random sample of scenarios.

Section 3.2 of the paper describes the problem and presents its several formulations in a mean-risk framework. Section 3.3 surveys related work on the problem, namely optimisation with risk measures and applications of metaheuristics to stochastic optimisation, portfolio selection and knapsack problems. The multiobjective local search metaheuristic is outlined in section 3.4, and section 3.5 presents the computational study. We conclude with a summary of the main contributions and future work perspectives in section 3.6.

3.2 The Static Stochastic Knapsack Problem with Random Weights

3.2.1 Problem Description

The Static Stochastic Knapsack Problem with Random Weights (SSKP-RW) can be described as the problem of choosing a subset of k items ($i = 1, \dots, k$), to be put into a knapsack of weight capacity q . Each item i has a reward r_i and a random weight W_i . Weight in excess is charged with a unit penalty c . The decision variables x_i take value 1 if item i is to be included in the solution (knapsack), and value 0 otherwise.

This problem has usually been defined considering the expected profit as the objective, thus leading to the following model:

$$\begin{aligned} \max \quad & \sum_{i=1}^k r_i x_i - cE \left[\max \left\{ \sum_{i=1}^k W_i x_i - q, 0 \right\} \right] \\ \text{s.t.} \quad & x_i \in \{0, 1\}, i = 1, \dots, k, \end{aligned} \tag{3.1}$$

where E denotes the expected value.

In the SSKP-RW, all items are simultaneously available, and the values of their weights are unknown before the inclusion decisions, that must be made concurrently. The problem has been subject to repeated studies mainly because it has several

practical applications and because many interesting stochastic optimisation problems have similar expected value objective functions (Kleywegt et al., 2002).

The SSKP-RW falls into a broader category of Stochastic Combinatorial Optimisation Problems (SCOP) with stochastic objective function that can be stated as follows:

$$\begin{aligned} \min \quad & E[f(\mathbf{x}, \omega)] \\ \text{s.t.} \quad & \mathbf{x} \in S, \end{aligned} \tag{3.2}$$

where \mathbf{x} is a solution for the problem, ω is the randomness component with a certain probability distribution, f is the objective function, E denotes the expected value, and S is the discrete, feasible region in the decision space.

Problems in this category are quite hard to tackle, due to their discrete nature and to the difficulties in evaluating, exactly or approximately, the objective function.

3.2.2 Mean-Risk Models

In contexts of decision making under risk, optimising a single expected value criterion will in general only be appropriate when the exact same decision situation occurs repeatedly, or when the decision maker is risk neutral. When these assumptions are not met, the inclusion of risk measures in stochastic models, leading to mean-risk models, provides an improved framework for decision support. Variance has classically been used as a risk measure, to a large extent due to Markowitz's influential work in portfolio management (Markowitz, 1959). In Markowitz's approach, two conflicting criteria are considered: the expected value of the portfolio's return, to be maximised; and the variance of the portfolio's return, to be minimised. This bicriteria optimisation problem can be solved by exploring the set of efficient solutions (those solutions for which improvement in one criterion is achieved only with the deterioration of the other) as a way to support the investor in expressing his implicit preferences and choosing a solution.

Research in risk measurement has pointed out several disadvantages in using variance as a risk measure, and put forward a number of alternatives to replace variance in the above formulation (cf. section 3.3). A measure that has been receiving increasing attention is the Conditional Value-at-Risk (CVaR_α), the conditional expected value beyond the Value-at-Risk (VaR_α). CVaR_α is a coherent risk measure and can be computed via linear programming (Rockafellar and Uryasev, 2002). Rigorous definitions of these measures require a certain number of subtleties to handle difficulties that may be presented by the probability functions. We refer, for instance, to Rockafellar and Uryasev (2002) for these rigorous definitions. In a simplified manner, considering a random variable representing profit, VaR_α could be defined as the minimum profit with probability level α and CVaR_α as the expected value of the profits below that minimum profit with probability level α .

Although mean-risk models have gained popularity in contexts of decision making under risk, their use has in general been limited to financial analysis. Only recently has the explicit consideration of risk concerns started to secure significant attention outside that domain, particularly as a result of the increasing research activity in stochastic optimisation and its applications. Some key references on risk measures and their use in Stochastic Integer Programming (SIP) may be found in Schultz (2003) and Ahmed (2006).

An application of this framework to SCOP might be described as a Mean-Risk Combinatorial Optimisation Problem, and stated as

$$\begin{aligned}
 \min \quad & \text{E}[f(\mathbf{x}, \omega)] \\
 \min \quad & \text{R}[f(\mathbf{x}, \omega)] \\
 \text{s.t.} \quad & \mathbf{x} \in S,
 \end{aligned} \tag{3.3}$$

where R is a risk function.

The general Mean-Risk SSKP-RW might accordingly be formulated in the follow-

ing way:

$$\begin{aligned}
& \max \quad \sum_{i=1}^k r_i x_i - cE \left[\max \left\{ \sum_{i=1}^k W_i x_i - q, 0 \right\} \right] \\
& \min \quad R \left[\sum_{i=1}^k r_i x_i - c \max \left\{ \sum_{i=1}^k W_i x_i - q, 0 \right\} \right] \\
& \text{s.t.} \quad x_i \in \{0, 1\}, i = 1, \dots, k.
\end{aligned} \tag{3.4}$$

Taking the variance as a risk measure, the risk objective function becomes:

$$\min \quad \text{Var} \left[\max \left\{ \sum_{i=1}^k W_i x_i - q, 0 \right\} \right], \tag{3.5}$$

where Var denotes the variance.

Additional difficulties arise for this kind of problems, from the multiobjective nature of the problem and from the quadratic nature of the variance objective. With CVaR_α as a risk measure, the risk objective function becomes the maximisation of the expected value of the profits below the minimum profit with probability level α , and can be stated as

$$\max \quad \sum_{i=1}^k r_i x_i - c\text{CVaR}_\alpha \left[\max \left\{ \sum_{i=1}^k W_i x_i - q, 0 \right\} \right], \tag{3.6}$$

where CVaR_α denotes the Conditional Value-at-Risk with probability level α .

These problems can be viewed as particular cases of Multiobjective Combinatorial Optimisation Problems, that can be represented by the following generic model:

$$\begin{aligned}
& \min \quad f_1(\mathbf{x}) = z_1 \\
& \quad \quad \quad \vdots \\
& \min \quad f_k(\mathbf{x}) = z_k \\
& \text{s.t.} \quad \mathbf{x} \in S,
\end{aligned} \tag{3.7}$$

where \mathbf{x} is a solution to the problem, S is the discrete, feasible region in the decision space, and f_1, \dots, f_k are the objective functions. $\mathbf{z} = (z_1, \dots, z_k)$ is called a criterion

vector. The feasible region in the objective space is $Z = \{\mathbf{z} \in \mathbb{R}^k : z_i = f_i(\mathbf{x}), \mathbf{x} \in S\}$. $\mathbf{z} \in Z$ is nondominated if and only if there is no other $\mathbf{z}' \in Z$ such that $z'_i \leq z_i, \forall i$, and $z'_i < z_i$, for some i . The nondominated set consists of all nondominated criterion vectors. $\mathbf{x} \in S$ is efficient if and only if its image in the objective space is nondominated. The efficient set consists of all efficient solutions.

As an important part of several methods for Multiobjective Combinatorial Optimisation, scalarising functions can be used for mapping criterion vectors to values in an ordinal scale of quality. The weighted sum scalarising function $s_{ws}(\mathbf{z}, \mathbf{z}^0, \lambda) = \sum_{i=1}^k \lambda_i (z_i - z_i^0)$, considers a reference criterion vector \mathbf{z}^0 and strictly positive scalar weights λ_i .

3.2.3 Sample Approximation

Difficulties in evaluating the expected value objective function, in an exact or in an approximate way, arise if a closed form does not exist, or if its values are hard to compute. These difficulties have often been handled with methods involving sample average approximation, where a random sample of N scenarios ω_j is generated so that the expected value function may be estimated from the sample average function. A discussion of the issues involved in this type of approach may be found in Kleywegt et al. (2002). This approximation may naturally be extended to the risk objective, and a *sample approximation for the mean-variance problem* would be formulated as

$$\begin{aligned} \min \quad & \frac{1}{N} \sum_{j=1}^N f(\mathbf{x}, \omega_j) \\ \min \quad & \frac{1}{N-1} \sum_{j=1}^N \left(f(\mathbf{x}, \omega_j) - \frac{1}{N} \sum_{j=1}^N f(\mathbf{x}, \omega_j) \right)^2 \\ \text{s.t.} \quad & \mathbf{x} \in S. \end{aligned} \tag{3.8}$$

A *sample approximation for the mean-CVaR $_\alpha$ problem* would be formulated as

(Rockafellar and Uryasev, 2002)

$$\begin{aligned}
\min \quad & \frac{1}{N} \sum_{j=1}^N f(\mathbf{x}, \omega_j) \\
\min \quad & \xi + \frac{1}{(1-\alpha)N} \sum_{j=1}^N Z_j \\
\text{s.t.} \quad & Z_j \geq f(\mathbf{x}, \omega_j) - \xi, \quad j = 1, \dots, N \\
& \mathbf{x} \in S, Z_j \geq 0, \xi \geq 0.
\end{aligned} \tag{3.9}$$

Considering weights W_i^j for each item i in each scenario j , the *sample approximation to the mean-variance SSKP-RW* is the following Quadratic Integer Programming problem:

$$\begin{aligned}
\max \quad & \sum_{i=1}^k r_i x_i - c \overline{Z}^+ \\
\min \quad & \frac{1}{N-1} \sum_{j=1}^N (Z_j^+ - \overline{Z}^+)^2 \\
\text{s.t.} \quad & Z_j^+ - Z_j^- = \sum_{i=1}^k W_i^j x_i - q, \quad j = 1, \dots, N \\
& Z_j^+ \leq \delta_j M, \quad j = 1, \dots, N \\
& Z_j^- \leq (1 - \delta_j) M, \quad j = 1, \dots, N \\
& \overline{Z}^+ = \frac{1}{N} \sum_{j=1}^N Z_j^+ \\
& x_i \in \{0, 1\}, \quad i = 1, \dots, k \\
& \delta_j \in \{0, 1\}, \quad j = 1, \dots, N \\
& Z_j^+, Z_j^- \geq 0, \quad j = 1, \dots, N \\
& \overline{Z}^+ \geq 0,
\end{aligned} \tag{3.10}$$

where $Z_j^+ - Z_j^-$ models the actual difference between the weight and the capacity in scenario j , δ_j are additional binary variables that take value 1 if the weight exceeds the capacity in scenario j or 0 otherwise, and M is an upper bound on the absolute values of the differences between total weight and capacity (M could be given, for example, by $\max_{j=1, \dots, N} \left\{ \sum_{i=1}^k W_i^j \right\}$). Z_j^+ and Z_j^- will have, for scenario j , the values of excess of weight and free capacity, respectively.

The *sample approximation to the mean-CVaR $_{\alpha}$ SSKP-RW* is an Integer Program-

ming problem with the following formulation:

$$\begin{aligned}
& \max \quad \sum_{i=1}^k r_i x_i - c \overline{Z}^+ \\
& \max \quad \sum_{i=1}^k r_i x_i - c \left(\xi + \frac{1}{(1-\alpha)N} \sum_{j=1}^N Y_j \right) \\
& \text{s.t.} \quad Z_j^+ \geq \sum_{i=1}^k W_i^j x_i - q, \quad j = 1, \dots, N \\
& \quad \quad \overline{Z}^+ = \frac{1}{N} \sum_{j=1}^N Z_j^+ \\
& \quad \quad Y_j \geq Z_j^+ - \xi, \quad j = 1, \dots, N \\
& \quad \quad Z_j^+, Y_j \geq 0, \quad j = 1, \dots, N \\
& \quad \quad \overline{Z}^+, \xi \geq 0.
\end{aligned} \tag{3.11}$$

Formulations (3.10) and (3.11) differ in the way that excess weight is modelled. The quadratic nature of the variance objective in the mean-variance formulation requires the use of additional binary variables δ_j to express the occurrence of overweight, whereas the mean-CVaR $_{\alpha}$ formulation handles the definition of excess weight through the interaction between objective functions and constraints.

3.2.4 Exact Objective Functions with Independent Normal Weights

If the weights of the items are independent and normally distributed, $W_i \sim N(\mu_i, \sigma_i^2)$, the random variable $Y(\mathbf{x}) = \sum_{i=1}^k W_i x_i - q$ is normally distributed with mean $\mu_Y(\mathbf{x}) = \sum_{i=1}^k \mu_i x_i - q$ and variance $\sigma_Y^2(\mathbf{x}) = \sum_{i=1}^k \sigma_i^2 x_i^2$. In this case the expected value, variance and CVaR $_{\alpha}$ of the excess weight $Z(\mathbf{x}) = \max\{Y(\mathbf{x}), 0\}$ can be computed exactly as

$$\begin{aligned}
\mathbb{E}[Z(\mathbf{x})] &= \sigma_Y(\mathbf{x}) \varphi\left(\frac{\mu_Y(\mathbf{x})}{\sigma_Y(\mathbf{x})}\right) + \mu_Y(\mathbf{x}) \Phi\left(\frac{\mu_Y(\mathbf{x})}{\sigma_Y(\mathbf{x})}\right) \\
\text{Var}[Z(\mathbf{x})] &= \mu_Y(\mathbf{x}) \sigma_Y(\mathbf{x}) \varphi\left(\frac{\mu_Y(\mathbf{x})}{\sigma_Y(\mathbf{x})}\right) + (\sigma_Y^2(\mathbf{x}) + \mu_Y^2(\mathbf{x})) \Phi\left(\frac{\mu_Y(\mathbf{x})}{\sigma_Y(\mathbf{x})}\right) - \mathbb{E}^2[Z(\mathbf{x})] \\
\text{CVaR}_\alpha[Z(\mathbf{x})] &= \begin{cases} \frac{\mathbb{E}[Z(\mathbf{x})]}{1-\alpha} & \text{if } \alpha \leq \Phi\left(-\frac{\mu_Y(\mathbf{x})}{\sigma_Y(\mathbf{x})}\right) \\ \mu_Y(\mathbf{x}) + \sigma_Y(\mathbf{x}) \frac{\varphi(\Phi^{-1}(\alpha))}{1-\alpha} & \text{if } \alpha > \Phi\left(-\frac{\mu_Y(\mathbf{x})}{\sigma_Y(\mathbf{x})}\right) \end{cases}
\end{aligned} \tag{3.12}$$

where φ denotes the standard normal probability density function and Φ denotes the standard normal probability distribution function.

3.3 Related Work

3.3.1 Static Stochastic Knapsack Problem

Alternative versions of the SSKP have been studied in the literature, by considering randomness in different subsets of the problem parameters. For problems with independent normally distributed rewards, Steinberg and Parks (1979) proposed a preference order dynamic programming algorithm, an approach further elaborated by Sniedovich (1980, 1981); Henig (1990) proposed a hybridisation of dynamic programming with a search procedure; Carraway et al. (1993) and Morton and Wood (1998) combined dynamic programming with branch-and-bound, for an objective of maximising the probability of a target achievement; Morton and Wood (1998) also developed a Monte Carlo approximation for problems with general distributions on the random rewards.

For problems with random weights and rewards that are linear functions of the weights, Cohn and Barnhart (1998) devised a branching approach, based on a binary tree, with items as nodes and inclusion decisions as branches.

For problems with random weights, Kleywegt et al. (2002) studied a Monte Carlo

simulation-based approach that repeatedly solves sample average optimisation problems, in which the expected value function is approximated by a sample average function, obtained by the generation of a random sample.

For problems with random capacity, Das and Ghosh (2003) proposed a depth first branch and bound algorithm and a heuristic based on local search using a two-swap neighbourhood structure.

The SSKP can be viewed as a Stochastic Integer Programming (SIP) problem, a broader class of problems for which there is extensive work reported in the literature, ranging from early work in heuristics to later developments with exact methods. Kleywegt et al. (2002) and Sahinidis (2004) present recent surveys of this area.

3.3.2 Optimisation with Risk Measures

The explicit treatment of risk in Stochastic Optimisation has only recently started to secure systematic attention. This is a field of active research, as it is also the case of the more specific issue of identifying adequate risk measures. For this adequacy, the prevailing concepts are consistency with stochastic dominance (Ogryczak and Ruszczyński, 1999) and coherence (Artzner et al., 1999). A common way to explicitly consider risk in stochastic optimisation problems is to include a second objective, in addition to expectation, consisting of a risk measure (such as dispersion parameters, excess probabilities, quantiles, or conditional expectations). Recent research in stochastic optimisation has focused on identifying measures that are coherent and consistent with stochastic dominance, while at the same time allowing the use of already available tools. Scalarisation approaches have therefore been privileged, with an emphasis in models with a weighted sum of the two objectives, often called “mean-risk” models. It should, however, be noted that in integer problems the feasible region may be nonconvex, and therefore nonsupported efficient solutions may exist which can not be found by optimising a weighted sum objective function.

Although mean-variance models have wide acceptance in practice, they are neither consistent with the stochastic dominance relation nor coherent (Ogryczak and Ruszczyński, 1999; Artzner et al., 1999). Additionally the quadratic nature of the variance objective does not allow an efficient use of mixed-integer linear programming solvers. These are two of the main reasons that have led researchers to consider approaches involving other risk measures.

Optimisation with conditional value-at-risk has been studied by Rockafellar and Uryasev (2002), Eichhorn and Romisch (2005) and Schultz and Tiedemann (2006). Schultz (2003) looks at multi-stage stochastic integer problems with excess probability, conditional value-at-risk and absolute semideviation. Ahmed (2006) investigates semideviation from a target, conditional value-at-risk, central semideviation, quantile-deviation and Gini's mean absolute difference. Shapiro and Ahmed (2004) examine a particular class of minimax stochastic programming models and relate it to mean-risk models with deviation from a quantile as risk measure. Takriti and Ahmed (2004) propose the use of non-decreasing variability measures, such as below fixed target risk measures, in the context of two-stage planning systems, to avoid suboptimality in the recourse problem. Kristoffersen (2005) considers central deviation, semideviation and expected excess of target. Markert and Schultz (2004) apply deviation based measures to SIP. Ruszczyński and Shapiro (2006) present a general theory of convex optimisation of convex risk measures.

3.3.3 Applications of Metaheuristics

Stochastic Optimisation Problems

The application of metaheuristics to stochastic optimisation problems has typically involved the incorporation of sampling methods for solution evaluation, and statistical inference methods for solution comparison. Alrefaie and Andradottir (1999), Ahmed

and Alkhamis (2002) and Rosen and Harmonosky (2005) are recent references on Simulated Annealing that provide overviews of developments in this field. Costa and Silver (1998) present an adaptation of Tabu Search along the generic lines mentioned above. In Haugen et al. (2001), scenario decomposition with the Progressive Hedging Algorithm (Rockafellar and Wets, 1991) is combined with Tabu Search to solve the mixed-integer sub-problems. In Gutjahr (2003), Ant Colony Optimisation uses Monte Carlo simulation for approximating the expected value objective function. Jin and Branke (2005) present an extensive survey on the use of Evolutionary Algorithms for optimisation in uncertain environments, devoting particular attention to multiobjective approaches, which allow the search for solutions with different tradeoffs between performance and robustness (Das, 2000; Ray, 2002; Jin and Sendhoff, 2003; Deb and Gupta, 2004). Cheng et al. (2004) also use a multiobjective evolutionary algorithm for multiobjective joint capacity planning and inventory control under uncertainty.

Portfolio Selection Problems

Mean-risk models have emerged and acquired relevance in the field of financial decision-making. It is also in this field that most applications of Multicriteria Decision Making (MCDM) techniques have been proposed for this type of models. Steuer and Na (2003) review applications of MCDM in finance, many of them in mean-risk models. Several algorithms based on metaheuristics have been proposed for non-linear mixed integer programming problems in portfolio selection. The addition of constraints on the number and proportion of assets in a portfolio, resulting in a mixed integer quadratic programming problem, is handled with Simulated Annealing, Tabu Search and Genetic Algorithms in Chang et al. (2000). In Crama and Schyns (2003) Simulated Annealing is used to approach another mixed integer quadratic programming formulation, that arises from the consideration of several practical constraints. Practical concerns, introducing non-linearity and integer valued variables, are again the

starting point for the work presented in Schlottmann and Seese (2004), in which a hybrid algorithm involving Multiobjective Evolutionary Algorithms and Local Search is used to approach a discrete risk-return efficient set, allowing for non-linear, non-quadratic, non-convex objective functions. In Ehrgott et al. (2004) a risk-return model with five objectives is proposed - a decision-maker utility function is built, based on a hierarchy of the objectives, and incorporated in a single objective nonlinear mixed integer programming model that is approached with Simulated Annealing, Tabu Search and Genetic Algorithms.

Knapsack Problems

In this context, metaheuristics have mostly been applied to a generalisation of the standard knapsack problem, called the multidimensional knapsack problem (MKP). The MKP extends the standard problem by considering several types of weights for each item, and a capacity constraint for each of these types of weights. In Freville (2004) a survey on this problem is presented, including a section on metaheuristics. The author mentions applications of Simulated Annealing, Tabu Search, Genetic Algorithms and Neural Networks, many of which make use of the properties of the problem to achieve improved results.

In this work we need to repeatedly solve instances of knapsack problems. However, as the focus is not on efficiency, we have adopted a rather straightforward algorithmic design and implementation. In this approach we have therefore directed our attention to the more basic components: solution representation, construction of initial solutions and local search neighbourhoods.

In the large majority of the literature, a standard binary string is used for solution representation (with the value 1 meaning the item is in the knapsack and the value 0 otherwise). This representation does not preclude infeasible solutions, and several alternative ways of dealing with infeasibility have been proposed. A review of these

approaches can be found in Hanafi and Freville (1998). In the particular problem we address, capacity constraints do not exist, thus leading to a problem formulation similar to the one proposed by Battiti and Tecchiolli (1992), who consider a penalty factor in the objective function to handle infeasibility.

Several simple and fast greedy algorithms for the MKP have been proposed in the literature, that can be used to provide initial solutions for metaheuristic algorithms. A section dedicated to reviewing greedy algorithms for the MKP is included in Freville (2004).

In local search based applications, the simplest movement that can be performed on a binary string is to change the value of a single item: an *add* movement will change it from 0 to 1; a *drop* movement from 1 to 0. More elaborate movements may be built from a set of strategically selected *drop* and *add* movements. Hanafi and Freville (1998) include a review of these movements.

Multiobjective metaheuristics have also been applied to multiobjective versions of knapsack problems, essentially for benchmarking purposes. In Teghem et al. (2000) an interactive procedure based on the author's multiobjective simulated annealing (MOSA) method is applied to a standard knapsack problem with multiple linear objectives. The same type of problem is handled in Gandibleux and Freville (2000) with a tabu search based procedure and decision space reduction. A survey and a benchmark of multiobjective evolutionary algorithms applied to another type of multiobjective knapsack problems (MOKP) are presented in Jaszkievicz (2002). This version of MOKP considers a set of items, a set of knapsacks and weights and profits, associating each item with each knapsack. For each knapsack, a capacity constraint is imposed and a profit function is to be maximised. More recently, in Phelps and Koksalan (2003) an interactive evolutionary algorithm has been proposed and applied to the standard knapsack problem with multiple linear objectives. Silva et al. (2006) present a scatter search based method for large size bi-criteria problems.

3.4 A Multiobjective Metaheuristic Approach

Multiobjective metaheuristics have been successfully applied to MOCO problems and are particularly well-suited to deal with the above mentioned difficulties arising in Mean-Risk Combinatorial Optimisation problems. Surveys on multiobjective metaheuristics are available in Ehrgott and Gandibleux (2000) and Jones et al. (2002).

Tabu Search for Multiobjective Combinatorial Optimisation (TAMOCO) (Hansen, 2000) and Pareto Simulated Annealing (PSA) (Czyzak and Jaskiewicz, 1998) can be viewed as Multiobjective Local Search (MOLS) approaches. Both aim at producing a good approximation of the efficient set, working with a population of solutions, each solution holding a weight vector for the definition of a search direction. Each approach proposes a different strategy for the definition of the weights, but share identical purposes for that definition: orientation of the search towards the nondominated frontier and spreading of solutions over that frontier (the former is achieved by the use of positive weights, while the latter is based on a comparison with other solutions of the population). Although in different ways, both methods operate on each single solution, searching and selecting a solution in its neighbourhood that will eventually replace it. Moreover, each procedure involves traditional metaheuristic components such as neighbourhoods, in general, or tabu lists, in the specific case of TAMOCO. The identification of these common aspects has suggested the definition of a MOLS generic template (Algorithm 5).

PSA and TAMOCO differ in the definition of several of the template's primitive operations: weight vectors are distinctly initialised and updated; $Neighbourhood(s)$ in PSA is a random subneighbourhood with just one movement; the generated movement in PSA is always selected, while in TAMOCO movement selection considers tabu status, aspiration criteria, and a comparison of evaluations based on a weighted sum scalarising function; in PSA a selected movement is accepted according to an

Algorithm 5: Multiobjective Local Search Template

```

Generate a set of initial feasible solutions  $G \subset S$ ;
Initialise the approximation to the efficient set  $E = \{\}$ ;
foreach  $s_i \in G$  do
    Initialise the corresponding context;
    Update  $E$  with  $s_i$ ;
end
while a stopping criterion is not met do
    foreach  $s_i \in G$  do
        Update the corresponding weight vector  $\lambda_i$ ;
        Initialise the selected solution  $s_s = 0$ ;
        foreach  $s' \in \text{Neighbourhood}(s_i)$  do
            Update  $E$  with  $s'$ ;
            if  $s'$  is selectable and  $s'$  is preferable to  $s_s$  then  $s_s = s'$ ;
        end
        if  $s_s \neq 0$  and  $s_s$  is acceptable then  $s_i = s_s$ ;
    end
end

```

acceptance probability, while in TAMOCO it is always accepted.

This template and related procedures have been implemented in an object-oriented framework called MetHOOD (Claro and Sousa, 2001) (Figure 3.1), that has been used to support the application described in this paper. Also of interest for this application is the support that the framework provides for neighbourhood variation, i.e., we can consider a sequence of neighbourhood structures and use them dinamically according to the evolution of the search process: if a new accepted solution is preferable to the current one, or if the current neighbourhood is the last in the sequence, the first neighbourhood in the sequence will be used next; otherwise the following neighbourhood in the sequence will be used next.

The MetHOOD framework has been instantiated for the SSKP-RW according to the following implementation choices:

- The solution representation is a binary string (where a 1 means the item is in the knapsack).
- For constructing initial solutions, the choice of items is based on a combina-

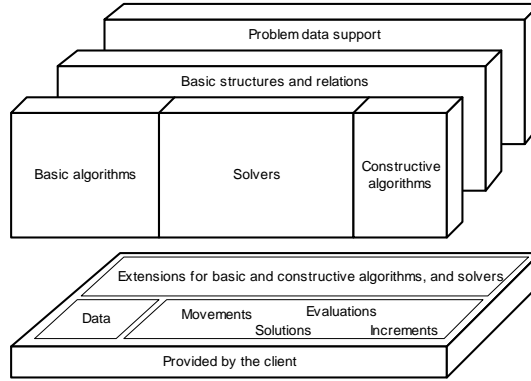


Figure 3.1: The MetHOOD framework

tion of one sorting criterion and one inclusion criterion: sorting criteria are *decreasing* r_i/μ_i or *decreasing* $r_i/(\mu_i + \sigma_i)$ ratios; inclusion criteria are *improving the expected value* or *keeping the probability of exceeding the capacity below a threshold value*.

- Neighbourhoods are built with one of the simplest movements for binary strings, the *flip* movement, which reverses the value of a binary string component.
- Objective functions for exact and approximate expected value, variance and CVaR_α have also been implemented. For CVaR_α , a value of $\alpha = 0.9$ has been considered. The normal distribution and its inverse were implemented with the Applied Statistics Algorithms AS66 (Hill, 1973) and AS241 (Wichura, 1988), respectively.

With this framework instantiation several MOLS algorithms become readily available. For computational experiments we have used an algorithm based in TAMOCO, without a tabu list, and using a variable neighbourhood.

3.5 Computational Study

3.5.1 Instances

Two sets of instances for the SSKP-RW have been generated, following Freville and Plateau (1994) and Kleywegt et al. (2002) (Table 3.1), one with a moderate size of 25 items and the other with a larger size of 250 items. Each of these sets includes 30 instances, 10 for each of 3 tightness factors. For each of these instances, one instance of an approximate problem with 1000 scenarios was generated. This number was chosen based on a computational study of the evolution of the approximation quality with the number of scenarios (as presented later in this section).

Table 3.1: Instance parameters

Parameter	Value
Number of items	25 and 250
Weights	Normal distribution
Weight mean	Uniform, between 50 and 100
Weight standard deviation	Uniform, between 5 and 10
Rewards	Mean weight added by a uniform value between 0 and 50
Unit penalty	5
Capacity	Sum of mean weights multiplied by a tightness factor
Tightness factor	0.25, 0.50 and 0.75

For the instances with 25 items, the nondominated sets could be obtained by full solution enumeration in short computational times (approximately 2 minutes for the exact problems, 10 minutes for the approximate problems with variance and 20 minutes for the approximate problems with CVaR_α). For the instances with 250 items, solution enumeration is not feasible anymore and our study focused on the approximate mean- CVaR_α instances, for which the nondominated sets could be obtained with an ϵ -constraint method, using ILOG CPLEX 10.1 MIP solver. For the approximate mean-variance instances, the QIP and QCIP solvers available in ILOG CPLEX 10.1 were used, but even for the instances with 25 items the computational times were quite large (hours).

All experiments were performed in a platform with an Intel Xeon Dual Core 5160 3.0 GHz CPU, 8 GB RAM, running Red Hat Enterprise Linux 4. The software was generated with GCC 3.4.6 with level 3 optimisation.

3.5.2 Performance Evaluation

The nondominated sets were used to evaluate the quality of the approximations. We have based our evaluation on one of the unary quality indicators with fewer limitations: the hypervolume (Zitzler and Thiele, 1998) bounded by the set $(\mathbf{z}^1, \mathbf{z}^2, \dots)$ and a reference point (\mathbf{z}^{ref}) (Figure 3.2). For each problem instance, a reference point has been chosen so that all points in the nondominated and approximation sets lie in the hypervolume, by considering the worst values for each objective function degraded by an additional 0.1%. A relative measure was built upon this one, consisting of the ratio between the values of the indicators for the approximation set and the nondominated set, so as to enable comparison of performance across multiple instances. The quality gap indicator being used consists of the difference to 1 of this measure.

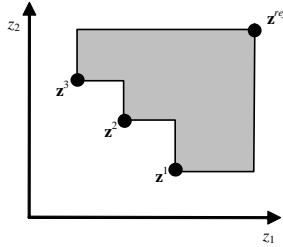


Figure 3.2: Hypervolume indicator

Considering, for example, the nondominated set $\{(1, 3), (2, 2), (3, 1)\}$ and an approximation set $\{(2, 4), (3, 3), (4, 2)\}$, for a problem with two minimisation objectives:

- the reference point is $(4.004, 4.004)$;
- the hypervolume of the nondominated set (shaded area in Figure 3.2) is 6.024016;

- the hypervolume of the approximation set is 1.016016;
- the quality gap indicator for the approximation set is $1 - 1.016016/6.024016 = 83.13\%$.

3.5.3 Study of Sample Approximation

To characterise the convergence behaviour of approximate problems, we have carried out a small study involving instances with 25 items. For each tightness factor, we have considered one exact problem instance and generated approximate problem instances with 50, 100, 200, 500, 1000 and 2000 scenarios. 10 instances were generated for each combination of tightness factor and number of scenarios. Through solution enumeration we have obtained the nondominated sets for all approximate problem instances. The solutions in these sets were then evaluated in the exact problem and two hypervolume indicators were recorded: one bounded by the nondominated solutions in the set (NDTD) and the other bounded by the nondominating solutions in the set (NDTG), providing information on how good the sets of, respectively, best and worst solutions are.

Average results are summarised in Figure 3.3 and in Table 3.2 and standard deviations in Table 3.3. Overall, we are able to verify the improvement in the quality of the approaches, with the nondominating sets converging to the nondominated sets, and both converging to the exact problem nondominated sets. This evolution is matched by the reduction in the standard deviations. It is also visible that for the same number of scenarios the approximation quality is lower for the mean-CVaR $_{\alpha}$ formulations, a fact that was expectable since CVaR $_{\alpha}$ is computed from the tail of the distribution. Relatively small improvements are obtained with the increase of the number of scenarios from 1000 to 2000.

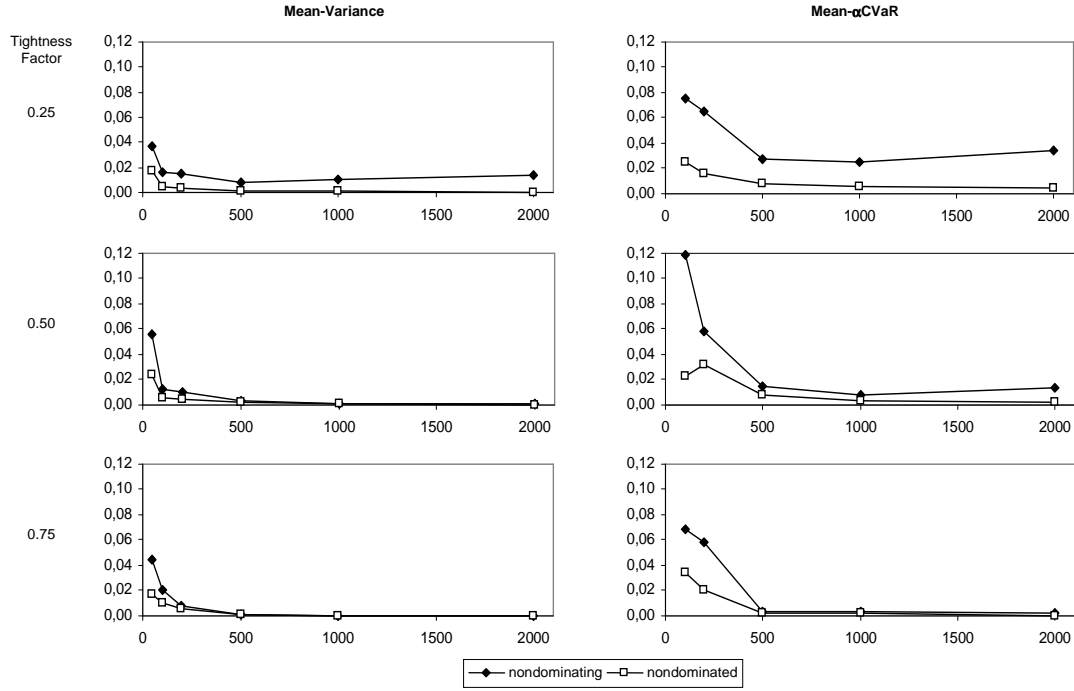


Figure 3.3: Average approximation quality gap as function of the number of scenarios

Table 3.2: Average approximation quality gap (%) as function of number of scenarios

Tight.	Set	Mean-Variance						Mean-CVaR $_{\alpha}$					
		50	100	200	500	1000	2000	50	100	200	500	1000	2000
0.25	NDTD	1.79	0.41	0.31	0.17	0.07	0.02	11.12	2.56	1.63	0.76	0.52	0.41
	NDTG	3.72	1.63	1.44	0.81	1.04	1.34	22.30	7.55	6.52	2.69	2.53	3.39
0.50	NDTD	2.35	0.56	0.50	0.18	0.08	0.04	5.50	2.29	3.22	0.79	0.33	0.24
	NDTG	5.61	1.24	1.00	0.35	0.17	0.08	25.97	11.92	5.79	1.51	0.75	1.41
0.75	NDTD	1.66	0.98	0.54	0.06	0.02	0.02	7.51	3.38	2.05	0.20	0.23	0.03
	NDTG	4.44	2.08	0.82	0.09	0.05	0.03	15.21	6.84	5.79	0.35	0.32	0.27

Table 3.3: Standard deviation of approximation quality gap (%) as function of number of scenarios

Tight.	Set	Mean-Variance						Mean-CVaR $_{\alpha}$					
		50	100	200	500	1000	2000	50	100	200	500	1000	2000
0.25	NDTD	1.60	0.18	0.14	0.09	0.07	0.01	12.24	1.87	0.87	0.72	0.72	0.24
	NDTG	3.95	1.15	0.85	0.76	0.76	0.64	24.60	6.74	7.97	1.35	1.41	1.35
0.50	NDTD	2.58	0.18	0.20	0.20	0.08	0.06	4.48	1.84	2.61	0.96	0.54	0.76
	NDTG	8.75	0.82	0.40	0.31	0.12	0.06	22.99	14.79	3.07	1.12	0.65	1.35
0.75	NDTD	0.58	0.56	0.53	0.04	0.03	0.02	4.29	2.48	2.58	0.34	0.34	0.10
	NDTG	4.04	2.43	0.52	0.04	0.02	0.02	13.09	7.01	5.37	0.40	0.39	0.37

3.5.4 Algorithm Configuration

The computational experiments have been performed with an adaptation of TAMOCO, as implemented in MethOOD, with no tabu list and with fixed or variable sub-neighbourhoods. These configurations can be viewed as a Multiobjective Random Local Search, in the case of fixed sub-neighbourhoods, and a Multiobjective Variable Neighbourhood Search, in the case of variable sub-neighbourhoods.

With the results of a series of preliminary algorithm executions we have confined the range of parameter values to be studied to those presented in Tables 3.4 and 3.5.

Each algorithm configuration has been executed 30 times for each instance. For all runs the generated approximation set has been recorded and its quality evaluated.

Table 3.4: Algorithm configurations for instances with 25 items

Parameter	Value
Sub-neighbourhood size	5, 10 and 20 movements
Variable neighbourhood	iterate sub-neighbourhoods of sizes 5, 10 and 20; with improvement return to size 5
Population	size 4 and 8 solutions
Constructive algorithms	equal number of solutions for each algorithm; excess weight threshold probability of 0.2
Time limit	1 second

Table 3.5: Algorithm configurations for instances with 250 items

Parameter	Value
Sub-neighbourhood size	100, 175 and 250 movements
Variable neighbourhood	iterate sub-neighbourhoods of sizes 100, 175 and 250; with improvement return to size 100
Population	size 16 and 32 solutions
Constructive algorithms	equal number of solutions for each algorithm; excess weight threshold probability of 0.2
Time limit	1 minute

3.5.5 Experimental Results

In tables 3.6, 3.7 and 3.8 we present the results obtained with the procedure configuration that provided higher quality results: for the problems with 25 items, the configuration with a population of 8 solutions and variable neighbourhood; for the problems with 250 items, the configuration with a population of 32 solutions and variable neighbourhood.

Overall, the results can be considered of high quality. For the exact problems with 25 items, the nondominated set is almost always found within the imposed time limit. The computational times for the mean-CVaR $_{\alpha}$ problem are lower than for the mean-variance problem. This is partially explained by the fact that the nondominated sets

Table 3.6: Results of computational study for exact instances with 25 items

Tight.	Inst.	Quality Gap (%)				Time (seconds)			
		Mean-Variance		Mean-CVaR $_{\alpha}$		Mean-Variance		Mean-CVaR $_{\alpha}$	
		Mean	St. Dev.	Mean	St. Dev.	Mean	St. Dev.	Mean	St. Dev.
0.25	1	0.00	0.00	0.00	0.00	0.70	0.20	0.02	0.01
	2	0.00	0.00	0.00	0.00	0.77	0.18	0.03	0.03
	3	0.00	0.00	0.00	0.00	0.75	0.20	0.00	0.01
	4	0.00	0.00	0.00	0.00	0.32	0.27	0.01	0.01
	5	0.00	0.00	0.00	0.00	0.64	0.20	0.01	0.01
	6	0.00	0.00	0.00	0.00	0.61	0.21	0.00	0.00
	7	0.00	0.00	0.00	0.00	0.78	0.19	0.01	0.01
	8	0.00	0.00	0.00	0.00	0.23	0.11	0.04	0.04
	9	0.00	0.00	0.00	0.00	0.67	0.22	0.00	0.01
	10	0.00	0.00	0.00	0.00	0.53	0.23	0.02	0.02
0.50	1	0.00	0.00	0.00	0.00	0.77	0.17	0.01	0.02
	2	0.00	0.00	0.00	0.00	0.87	0.09	0.22	0.22
	3	0.00	0.00	0.00	0.00	0.88	0.13	0.12	0.10
	4	0.00	0.00	0.00	0.00	0.85	0.13	0.10	0.11
	5	0.00	0.00	0.00	0.00	0.87	0.11	0.16	0.09
	6	0.00	0.00	0.00	0.00	0.87	0.11	0.04	0.07
	7	0.01	0.03	0.00	0.00	0.85	0.15	0.06	0.04
	8	0.01	0.02	0.00	0.00	0.92	0.07	0.55	0.30
	9	0.02	0.05	0.01	0.02	0.95	0.04	0.54	0.26
	10	0.00	0.00	0.00	0.00	0.65	0.26	0.04	0.04
0.75	1	0.00	0.02	0.08	0.42	0.88	0.11	0.27	0.21
	2	0.00	0.00	0.00	0.00	0.84	0.17	0.08	0.06
	3	0.01	0.03	0.37	1.48	0.87	0.14	0.18	0.26
	4	0.00	0.00	0.00	0.00	0.89	0.10	0.10	0.09
	5	0.00	0.00	0.00	0.00	0.85	0.11	0.29	0.18
	6	0.00	0.00	0.00	0.00	0.86	0.09	0.02	0.02
	7	0.00	0.00	0.00	0.00	0.85	0.14	0.04	0.04
	8	0.00	0.01	0.00	0.00	0.90	0.06	0.41	0.26
	9	0.04	0.01	0.00	0.00	0.87	0.12	0.53	0.25
	10	0.00	0.00	0.00	0.00	0.85	0.13	0.01	0.01

for the former have less solutions than for the latter.

For the sample problems with 25 items, the approximation quality is again very high, with difficulties arising in just 4 instances, for which the average quality gap indicator values remain above 2%. The computational times for the mean-CVaR $_{\alpha}$ problem are closer to the times for the mean-variance problem, due to the higher computational effort required to compute CVaR $_{\alpha}$. They are, still, at least an order of magnitude lower than the times required by the ϵ -constraint method.

For the sample problems with 250 items, the average quality gap indicator values are below 1% for 14 instances, between 1% and 2% for 8 instances, and above 2% for the remaining 8 instances, with a higher value of 4.34%. These results were achieved

Table 3.7: Results of computational study for approximate instances with 25 items

Tight.	Inst.	Quality Gap (%)				Time (seconds)				
		Mean-Variance		Mean-CVaR $_{\alpha}$		Mean-Variance		Mean-CVaR $_{\alpha}$		ϵ -constraint
		Mean	St. Dev.	Mean	St. Dev.	Mean	St. Dev.	Mean	St. Dev.	
0.25	1	0.11	0.21	0.00	0.00	0.69	0.20	0.55	0.27	29.89
	2	0.99	1.50	0.24	1.00	0.60	0.28	0.48	0.26	16.55
	3	1.20	1.71	0.00	0.00	0.72	0.20	0.10	0.07	12.91
	4	0.93	1.17	0.24	0.74	0.72	0.23	0.34	0.22	17.58
	5	0.02	0.05	0.02	0.05	0.83	0.14	0.57	0.25	22.78
	6	0.00	0.01	0.00	0.00	0.57	0.28	0.06	0.06	5.51
	7	0.00	0.01	0.00	0.01	0.48	0.22	0.15	0.17	7.28
	8	0.02	0.03	0.00	0.00	0.88	0.10	0.58	0.20	32.52
	9	0.04	0.09	0.01	0.03	0.65	0.26	0.21	0.22	6.49
	10	0.00	0.00	0.00	0.00	0.74	0.18	0.18	0.17	16.60
0.50	1	0.06	0.18	0.03	0.08	0.56	0.27	0.37	0.22	9.09
	2	0.24	0.09	4.93	9.25	0.80	0.16	0.38	0.23	14.40
	3	2.34	1.14	2.94	2.61	0.72	0.23	0.39	0.26	6.35
	4	0.80	0.59	1.21	1.30	0.78	0.17	0.45	0.24	19.28
	5	0.22	0.19	0.15	0.19	0.84	0.12	0.69	0.22	22.96
	6	0.05	0.10	0.00	0.00	0.67	0.22	0.21	0.24	8.88
	7	0.03	0.03	0.01	0.03	0.62	0.21	0.66	0.22	19.46
	8	0.09	0.06	0.02	0.07	0.83	0.17	0.30	0.23	24.66
	9	0.21	0.14	0.31	0.25	0.79	0.14	0.66	0.26	30.36
	10	0.00	0.00	0.00	0.00	0.75	0.18	0.61	0.22	14.26
0.75	1	2.79	2.60	5.95	9.11	0.85	0.13	0.58	0.27	9.37
	2	0.03	0.04	0.06	0.07	0.68	0.26	0.56	0.27	12.30
	3	0.51	1.12	0.30	0.55	0.73	0.19	0.17	0.17	10.01
	4	0.76	1.01	2.81	3.29	0.88	0.12	0.51	0.27	12.97
	5	0.15	0.07	0.05	0.07	0.58	0.22	0.37	0.30	14.71
	6	0.01	0.01	0.00	0.02	0.79	0.14	0.27	0.22	9.52
	7	0.03	0.08	0.00	0.00	0.69	0.26	0.42	0.24	9.76
	8	0.01	0.00	0.11	0.08	0.79	0.18	0.41	0.26	18.41
	9	0.14	0.05	0.01	0.00	0.72	0.19	0.46	0.32	9.78
	10	0.00	0.00	0.00	0.00	0.80	0.13	0.44	0.22	9.13

Table 3.8: Results of computational study for approximate instances with 250 items

Tight.	Inst.	Mean-CVaR $_{\alpha}$				
		Quality Gap (%)		Time (seconds)		
		Mean	St. Dev.	Mean	St. Dev.	ϵ -constraint
0.25	1	1.13	0.27	58.07	3.52	1891.70
	2	2.37	0.31	56.79	4.38	2471.28
	3	2.48	0.47	55.88	5.65	2768.80
	4	3.51	0.69	59.52	1.81	3327.54
	5	4.41	0.73	59.41	1.15	1698.48
	6	1.04	0.30	56.50	4.30	1480.96
	7	4.34	0.95	59.54	1.61	2180.45
	8	2.46	0.36	58.39	2.32	3028.53
	9	1.15	0.24	58.30	2.72	4368.35
	10	2.23	0.64	56.69	3.97	1410.78
0.50	1	0.61	0.26	59.25	1.72	878.41
	2	0.51	0.22	59.70	1.37	1204.14
	3	1.41	0.20	59.57	1.73	1245.36
	4	1.15	0.23	59.60	1.61	1209.51
	5	0.12	0.08	56.03	4.94	513.06
	6	0.99	0.26	58.98	2.32	1108.42
	7	1.71	0.50	59.79	0.84	1372.04
	8	1.08	0.26	58.19	2.64	1127.13
	9	2.52	0.53	59.95	1.04	1996.58
	10	0.48	0.16	57.23	3.55	503.37
0.75	1	1.44	0.42	59.43	1.36	1626.50
	2	0.82	0.17	60.02	1.10	1268.57
	3	0.90	0.20	59.84	1.14	2097.67
	4	0.82	0.27	59.15	1.59	1347.07
	5	0.51	0.18	58.72	2.76	1009.84
	6	0.56	0.13	59.99	0.96	1020.38
	7	0.36	0.10	59.66	1.33	1142.51
	8	0.20	0.07	59.28	1.38	825.04
	9	0.33	0.11	58.85	2.53	659.51
	10	0.39	0.09	57.98	2.81	711.31

with a computational time limit that is again at least an order of magnitude lower than the times required by the ϵ -constraint method.

Figure 3.4 presents approximation sets with values for the quality gap indicator of 15%, 10%, 5% and 2%, and compares them to the corresponding nondominated set, to provide a more accurate notion of the approximation quality for several levels of the quality gap indicator.

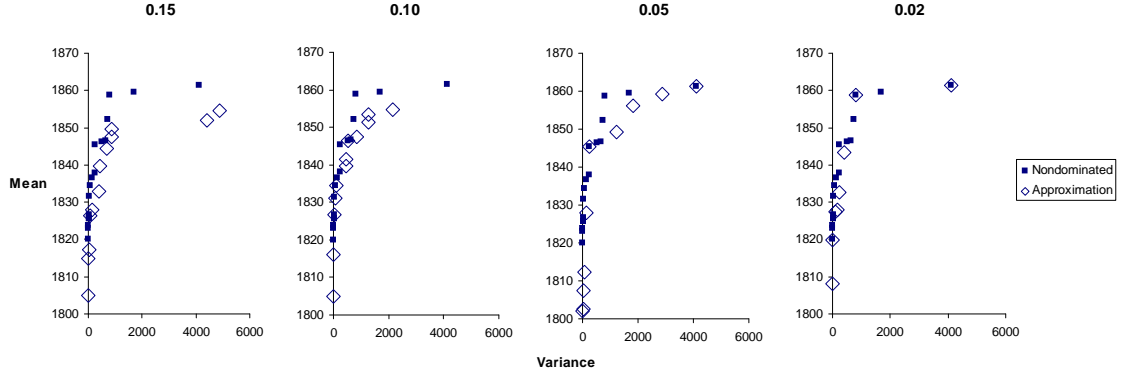


Figure 3.4: Nondominated sets and approximation sets of different quality gap levels

3.6 Conclusions

The work described in this paper goes beyond what has been reported in the literature for the SSKP, by introducing an approach that considers mean and risk criteria, and can handle both exact and sample approximation problems. IP and QIP/QCIP solvers are unable to tackle the exact problem, but can be used to obtain efficient sets for the sample approximation problems, for example using the ϵ -constraint method. However, the computational times involved are very large, whereas the approach presented here can produce high quality approximations to the efficient sets in very short computational times.

Another unique feature of this approach is the fact that adopting different risk measures can be done by simply changing the corresponding objective function, while keeping the remaining parts of the implementation. Multiobjective metaheuristics have had very confined applications in optimisation with uncertainty, in particular in the areas of robust optimisation and portfolio selection. With this work we explicitly introduce a multiobjective mean-risk framework for the general class of Stochastic Combinatorial Optimisation problems and show that multiobjective metaheuristics

are a class of algorithms that are well-suited to deal with the difficulties presented by these problems.

This work is being followed by an effort to apply the mean-risk tools and framework to the area of Operations Strategy, where we are particularly interested in problems of capacity investment in contexts of uncertainty, that have seldom been approached explicitly considering risk.

Chapter 4

A Multiobjective Metaheuristic for a Mean-Risk Multistage Capacity Investment Problem

(Under review at Journal of Heuristics)

In this paper, we propose a multiobjective local search metaheuristic for a mean-risk multistage capacity investment problem with irreversibility, lumpiness and economies of scale in capacity costs. Conditional value-at-risk is considered as a risk measure. Results of a computational study are presented, that clearly indicate that the approach is capable of producing high-quality approximations to the efficient sets, with a modest computational effort.

4.1 Introduction

Decisions about capacity investments are among the most important in the context of Operations Strategy. They are basically concerned with the choice of type, magnitude,

timing and location of capacity investments. These decisions are often irreversible and their impact can be extremely important. In fact, a large mismatch, either positive or negative, between capacity and demand will always be negatively reflected on the performance of a company.

Uncertainty in demand and in the conditions of resource availability is intrinsic to the long-term settings of decisions concerning capacity investments, and naturally increases the complexity of these decisions. Van Mieghem (2003) points out that the consideration of the variability of payoffs in uncertain settings has received very scarce attention in the literature. This complexity can be even higher in the presence of frictions such as lumpiness or economies of scale in capacity costs, which may result from technological factors or from the configuration of the suppliers' operations. Wu et al. (2005) recognise the presence of the above-mentioned features in several industries and the difficulty in dealing with them, particularly in high-tech sectors (semiconductors, consumer electronics, telecommunications, and pharmaceutical). In fact, Van Mieghem (2003) identifies indivisibility, irreversibility and nonconvexity as the three main challenges in capacity cost modelling.

In this paper we address a problem of multistage capacity investment under uncertainty, with irreversibility, lumpiness and economies of scale in capacity costs. The uncertainty in demand and costs is considered through a scenario tree. The economies of scale are modelled with fixed charge cost functions. The problem is approached from a mean-risk perspective, considering conditional value-at-risk (CVaR) as a risk measure, with an application of multiobjective metaheuristics to a bi-objective mean-CVaR formulation.

The paper is organised as follows: section 4.2 briefly reviews the relevant literature; section 4.3 presents an integer programming formulation for the problem; section 4.4 describes the multiobjective metaheuristics approach; section 4.5 presents the computational experiments and results; section 4.6 closes the paper with conclusions

and perspectives for future work.

4.2 Related Work

The work presented in this paper has its starting point in the problem of multistage capacity expansion under uncertainty considered in Ahmed et al. (2003) and Ahmed and Sahinidis (2003). These authors use a scenario tree to model the stochastic evolution of costs and demand, and consider economies of scale in costs through fixed-charge cost functions. The resulting problem is a multistage stochastic integer programming problem, a class of problems for which the stochastic programming literature provides no generic approaches. The authors develop approximation and reformulation schemes that lead to significant improvements on the computational times obtained by straightforward use of IP solvers.

The survey on capacity management presented in Van Mieghem (2003) has motivated us to readdress this problem, by including some additional challenging features in capacity modelling, namely lumpiness and the explicit consideration of risk. This motivation has been recently reinforced by the survey on capacity management in high-tech industries by Wu et al. (2005), that confirms the relevance of these features for capacity investment decisions in those industries. These recent references have presented comprehensive reviews of literature, and therefore we direct interested readers to them for in depth information on the theory and applications of capacity management and, in particular, on issues related to capacity expansion under uncertainty.

Recent advances in the explicit consideration of risk in stochastic optimisation have been stimulated by advances in the identification of adequate risk measures. Conditional value-at-risk (CVaR_α), the conditional expected value beyond value-at-risk (VaR_α), is a risk measure that has received significant attention, mainly due to

the fact that it is coherent and can be computed via linear programming (Rockafellar and Uryasev, 2002). Rigorous definitions of these measures require a certain number of subtleties to handle the difficulties raised by the probability functions. We refer, for instance, to Rockafellar and Uryasev (2002) for these rigorous definitions. In a simplified way, given a random variable representing cost, VaR_α could be defined as the maximum cost with probability level α and CVaR_α as the expected value of the costs above that maximum cost with probability level α .

Optimisation with conditional value-at-risk has also been studied by Eichhorn and Romisch (2005) and Schultz and Tiedemann (2006). Schultz (2003) looks at multistage stochastic integer problems with excess probability, conditional value-at-risk and absolute semideviation. Ahmed (2006) investigates semideviation from a target, conditional value-at-risk, central semideviation, quantile-deviation and Gini's mean absolute difference. Ruszczyński and Shapiro (2006) present a general theory of convex optimisation of convex risk measures. The emphasis in a great majority of these works has been on models that consider a weighted sum of the mean and risk objectives. In integer problems this type of approaches has the disadvantage of being unable to find nonsupported efficient solutions.

The application of metaheuristics to stochastic optimisation problems has typically involved the incorporation of sampling methods for solution evaluation, and statistical inference methods for solution comparison. Alrefaei and Andradottir (1999), Ahmed and Alkhamis (2002) and Rosen and Harmonosky (2005) are recent references on Simulated Annealing that provide overviews of developments in this field. Costa and Silver (1998) present an adaptation of Tabu Search along the generic lines mentioned above. In Haugen et al. (2001), scenario decomposition with the Progressive Hedging Algorithm (Rockafellar and Wets, 1991) is combined with Tabu Search to solve the mixed-integer sub-problems. In Gutjahr (2003), Ant Colony Optimisation uses Monte Carlo simulation for approximating the expected value objective func-

tion. Jin and Branke (2005) present an extensive survey on the use of Evolutionary Algorithms for optimisation in uncertain environments, devoting particular attention to multiobjective approaches that allow the search for solutions with different trade-offs between performance and robustness (Das, 2000; Ray, 2002; Jin and Sendhoff, 2003; Deb and Gupta, 2004). Cheng et al. (2004) also use a multiobjective evolutionary algorithm for multiobjective joint capacity planning and inventory control under uncertainty.

Several algorithms based on metaheuristics have been proposed for risk-return non-linear MIP problems in portfolio selection. Quadratic formulations with cardinality/proportion and other practical constraints have been addressed in Chang et al. (2000) with Simulated Annealing, Tabu Search and Genetic Algorithms, and in Crama and Schyns (2003) with Simulated Annealing. Practical concerns, reflected in the introduction of non-linearity and integer valued variables, are again the starting point for the work presented in Schlottmann and Seese (2004) with a multiobjective hybridisation of evolutionary and local search algorithms. In Ehrgott et al. (2004) a risk-return model with five objectives is proposed and approached with Simulated Annealing, Tabu Search and Genetic Algorithms.

4.3 A Mean-Risk Model

We consider here a discretised planning horizon, over which the evolution of demand and costs is modelled with a scenario tree. Each level of the tree corresponds to a time period. $\mathcal{T}(n)$ denotes the subtree rooted in node n , with $n = 0$ being the root node, and $\mathcal{P}(n)$ the path from the root node to node n . \mathcal{S} is the set of leaf nodes, each corresponding to one of the S equally probable scenarios. Each node $n \in \mathcal{T}(0)$ is characterised by demand d_n . In each node n , each resource $i \in \mathcal{I}$ is characterised by a discounted fixed cost $\alpha_{i,n}$ and a discounted variable cost $\beta_{i,n}$ of expansion. The

capacity of resource i can be changed by discrete increments of value l_i . The supply of capacity is unbounded and initial capacities are zero, without loss of generality. The decision variables are $x_{i,n}$, the number of capacity increments for resource i at node n . The binary variables $y_{i,n}$ take the value 1 if capacity of resource i is incremented in node n , and the value 0 otherwise.

Figures 4.1 and 4.2 show a binary scenario tree for a problem with 3 resources ($i = 1, 2, 3$) and 3 periods, and a corresponding feasible solution:

- Figure 4.1 shows the demand in each node satisfied by the sum of all capacity expansions in the path from the root to that node. In each node, from left to right, the first column contains the capacity expansions performed for each resource in that node, the second shows the capacity for each resource, with total capacity in the bottom row, and the third, with just one row, is the demand.

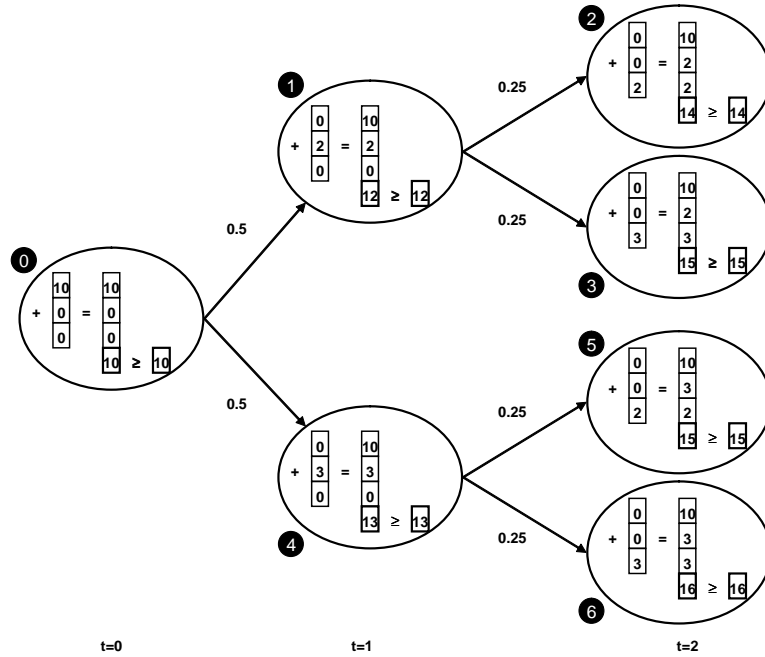


Figure 4.1: Solution and demand constraints

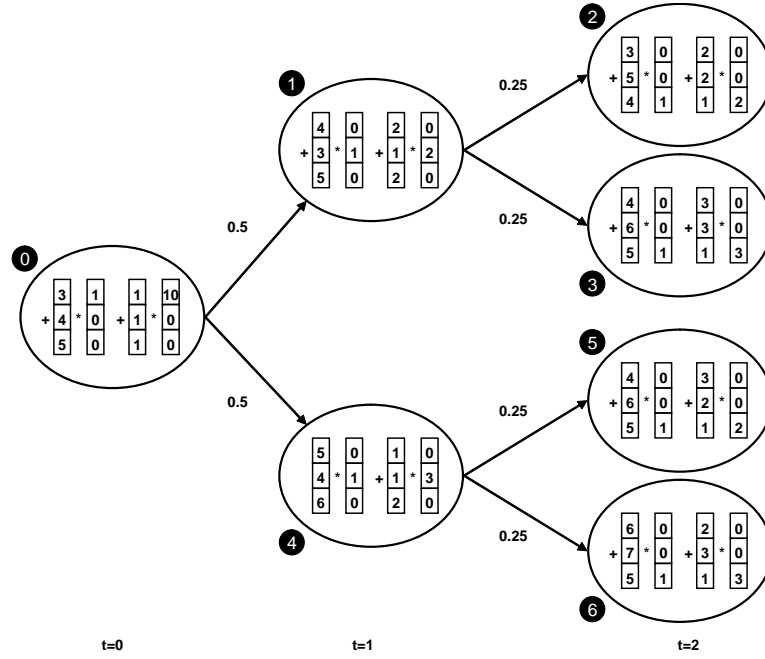


Figure 4.2: Solution and cost structure

- Figure 4.2 displays the cost structure for the same solution, with the two components considered for each resource in each node - fixed and variable costs. In each node, from left to right, the columns contain: the fixed costs, the binary variables $y_{i,n}$, the variable costs and the capacity expansions.

In both figures the solution is represented by the capacity expansions in each node. To each arc is associated the probability of the node where the arc is directed to.

The inclusion of a CVaR_α risk objective results in the following bi-objective integer programming formulation:

$$\begin{aligned}
\min \quad & E = \frac{1}{S} \sum_{m \in \mathcal{S}} \sum_{n \in \mathcal{P}(m)} \sum_{i \in \mathcal{I}} (\alpha_{i,n} y_{i,n} + \beta_{i,n} l_i x_{i,n}) \\
\min \quad & \text{CVaR}_\alpha = \xi + \frac{1}{(1-\alpha)S} \sum_{m \in \mathcal{S}} Z_m \\
\text{s.t.} \quad & Z_m \geq \sum_{n \in \mathcal{P}(m)} \sum_{i \in \mathcal{I}} (\alpha_{i,n} y_{i,n} + \beta_{i,n} l_i x_{i,n}) - \xi, \quad m \in \mathcal{S} \\
& \sum_{p \in \mathcal{P}(n)} \sum_{i \in \mathcal{I}} l_i x_{i,p} \geq d_n, \quad n \in \mathcal{T}(0) \\
& l_i x_{i,n} \leq M y_{i,n}, \quad n \in \mathcal{T}(0), i \in \mathcal{I} \\
& x_{i,n} \geq 0 \text{ and integer}, \quad n \in \mathcal{T}(0), i \in \mathcal{I} \\
& y_{i,n} \in \{0, 1\}, \quad n \in \mathcal{T}(0), i \in \mathcal{I} \\
& Z_m \geq 0, \quad m \in \mathcal{S} \\
& \xi \geq 0.
\end{aligned} \tag{4.1}$$

In this formulation, the objective E minimises the *expected value of the total investment cost* over the planning horizon, and the objective CVaR_α minimises the *conditional value-at-risk of that investment cost*. The first set of constraints are the usual constraints required for defining the conditional value-at-risk. The second set of constraints are the demand satisfaction constraints. The third set of constraints define $y_{i,n}$ in terms of variables $x_{i,n}$.

4.4 A Multiobjective Metaheuristic Approach

Multiobjective metaheuristics have been successfully applied to Multiobjective Combinatorial Optimisation (MOCO) problems and are particularly well-suited to deal with the set of previously mentioned challenging features in capacity modelling. Surveys on multiobjective metaheuristics are available in Ehrgott and Gandibleux (2000) and Jones et al. (2002).

Tabu Search for Multiobjective Combinatorial Optimisation (TAMOCO) (Hansen, 2000) and Pareto Simulated Annealing (PSA) (Czyzak and Jaskiewicz, 1998) can be viewed as Multiobjective Local Search (MOLS) approaches. Both aim at producing a good approximation of the efficient set, working with a population of solutions, each solution holding a weight vector for the definition of a search direction. Each approach proposes a different strategy for the definition of the weights, but with the same purpose: orientation of the search towards the nondominated frontier and spreading of solutions over that frontier (the former is achieved by the use of positive weights, while the latter is based on a comparison with other solutions of the population). Although in different ways, both methods operate on each single solution, searching and selecting a solution in its neighbourhood that will eventually replace it. Moreover, each procedure involves traditional metaheuristic components such as neighbourhoods, in general, or tabu lists, in the specific case of TAMOCO. The identification of these common aspects has suggested the definition of a MOLS generic template (Algorithm 6).

Algorithm 6: Multiobjective Local Search Template

```

Generate a set of initial feasible solutions  $G \subset S$ ;
Initialise the approximation to the efficient set  $E = \{\}$ ;
foreach  $s_i \in G$  do
    Initialise the corresponding context;
    Update  $E$  with  $s_i$ ;
end
while a stopping criterion is not met do
    foreach  $s_i \in G$  do
        Update the corresponding weight vector  $\lambda_i$ ;
        Initialise the selected solution  $s_s = 0$ ;
        foreach  $s' \in \text{Neighbourhood}(s_i)$  do
            Update  $E$  with  $s'$ ;
            if  $s'$  is selectable and  $s'$  is preferable to  $s_s$  then  $s_s = s'$ ;
        end
        if  $s_s \neq 0$  and  $s_s$  is acceptable then  $s_i = s_s$ ;
    end
end

```

PSA and TAMOCO differ in the definition of several of the template's primitive operations: weight vectors are distinctly initialised and updated; *Neighbourhood*(s) in PSA is a random subneighbourhood with just one movement; the generated movement in PSA is always selected, while in TAMOCO movement selection considers tabu status, aspiration criteria, and a comparison of evaluations based on a weighted sum scalarising function; in PSA a selected movement is accepted according to an acceptance probability, while in TAMOCO it is always accepted.

This template and related procedures have been implemented in an object-oriented framework called MetHOOD (Claro and Sousa, 2001) (Figure 4.3), that has been used to support the application described in this paper. Also of interest for this application is the support that the framework provides for neighbourhood variation, i.e., we can consider a sequence of neighbourhood structures and use them dynamically according to the evolution of the search process: if a new accepted solution is preferable to the current one, or if the current neighbourhood is the last in the sequence, the first neighbourhood in the sequence will be used next; otherwise the following neighbourhood in the sequence will be used next.

The MetHOOD framework has been instantiated for the capacity investment prob-

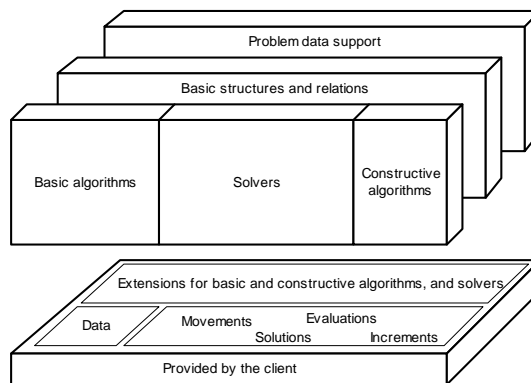


Figure 4.3: The MetHOOD framework

lem according to the implementation choices described next. With this framework instantiation several MOLS algorithms become readily available.

4.4.1 Solution

A tree representation has been used, matching the scenario tree structure that models demand and costs (see Section 4.3), with an array of capacity expansions for each resource in each node. Any node with a maximum demand in its subtree not higher than the maximum demand in the path from the root to its parent is not considered in the solution representation, since existing capacity at that node will be able to satisfy all the demand in its subtree with no further expansions.

4.4.2 Capacity Variation (CV) Neighbourhood

This neighbourhood structure is defined by modifications to a solution consisting of positive or negative capacity variations, for a selected resource in a selected node.

Positive Variations

Positive variations are limited by the maximum demand in the subtree of the selected node. The nodes in this subtree are visited in depth first order: if the capacity in a node does not exceed the new capacity value by more than the node's *capacity surplus* (the minimum difference between capacity and demand at any node in its subtree), any existing expansions are discarded; otherwise, existing expansions are reduced to the required level, starting with the resources that provide higher cost reduction, and the nodes below are not visited, their capacity expansions remaining unchanged.

Figures 4.4 and 4.5 illustrate a positive capacity variation for resource 2 in node 0, from 0 to 2, with total capacity increased from 10 to 12. The considered costs are the ones presented in Figure 4.2. The boxes with black background correspond to the

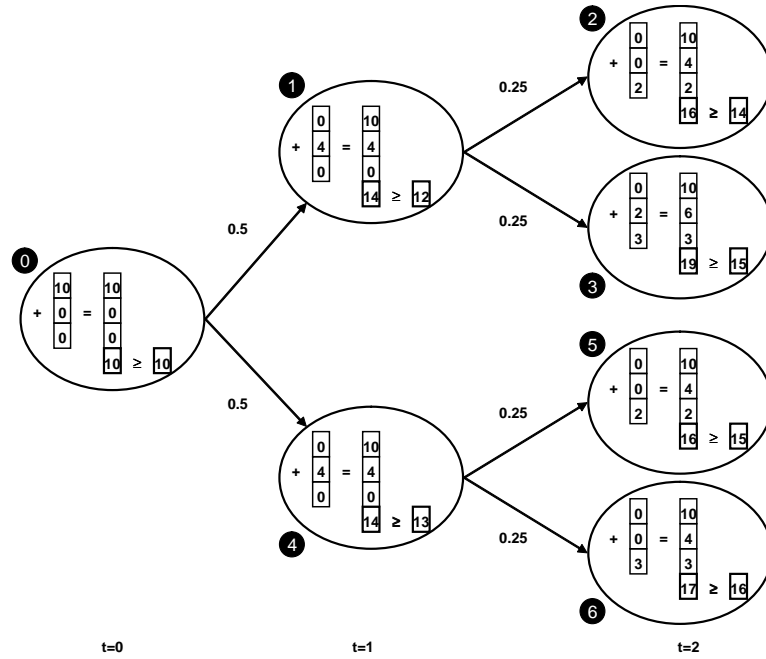


Figure 4.4: Solution before positive capacity variation

adjustments in capacity expansions:

1. In node 1, capacity was 14 with a surplus of 2. The difference to 12 does not exceed the surplus, so all expansions are discarded.
2. In node 2, capacity was 16 with a surplus of 2. The difference to 12 exceeds the surplus, but the minimum required capacity is 14, so expansions remain unchanged.
3. In node 3, capacity was 19 with a surplus of 4. The difference to 12 exceeds the surplus, so expansions are adjusted. The minimum required capacity is 15, so a total of 2 may be discarded. The resource where the cost savings with this reduction are greater is resource 2, so its expansion is completely discarded.
4. In node 4, capacity was 14 with a surplus of 1. The difference to 12 exceeds the

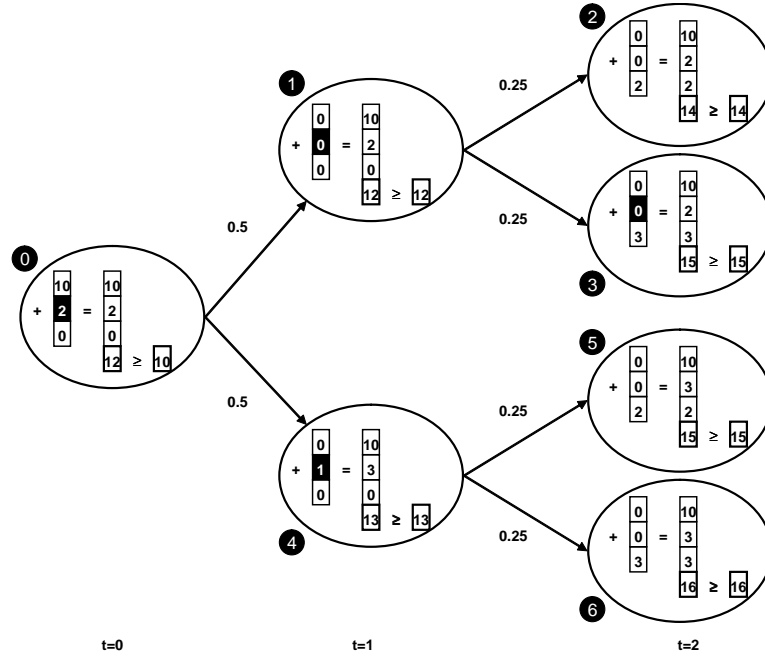


Figure 4.5: Solution after positive capacity variation

surplus, so expansions are adjusted. The minimum required capacity is 13, so a total of 3 may be discarded. The only resource with an expansion is resource 2, so its expansion is reduced to 1. The expansions in the node's subtree remain unaltered.

Negative Variations

Negative variations are limited by the capacity of the selected resource in the selected node, and by the demand satisfaction constraints for the nodes in the path from the selected node to the root. The capacity level has to be adjusted for the set of nodes in the path from the selected node to the root where the selected resource's capacity level is higher or equal to the reduced value: in the topmost of these nodes, the resource's expansion is reduced so that its capacity becomes equal to the reduced value; in the remaining nodes, existing expansions of the selected resource are discarded. In

the children of the nodes in this set that are not themselves part of the set, an expansion covering the capacity reduction in the parent node, and considering the node's capacity surplus, is created for the resource that provides it at the minimum cost.

Figures 4.6 and 4.7 illustrate a negative capacity variation for resource 1 in node 2, from 12 to 10. Costs are presented in Figure 4.2. The boxes with black background correspond to the adjustments in capacity expansions:

1. The capacity reduction for resource 1 in node 2 is limited by the capacity and demand conditions in nodes 2, 1 and 0. The minimum required capacities for resource 1, in these nodes, are 8, 6 and 10, respectively. The capacity can not, therefore, be reduced below 10, so the maximum feasible reduction is 2. The proposed reduction is within this range.

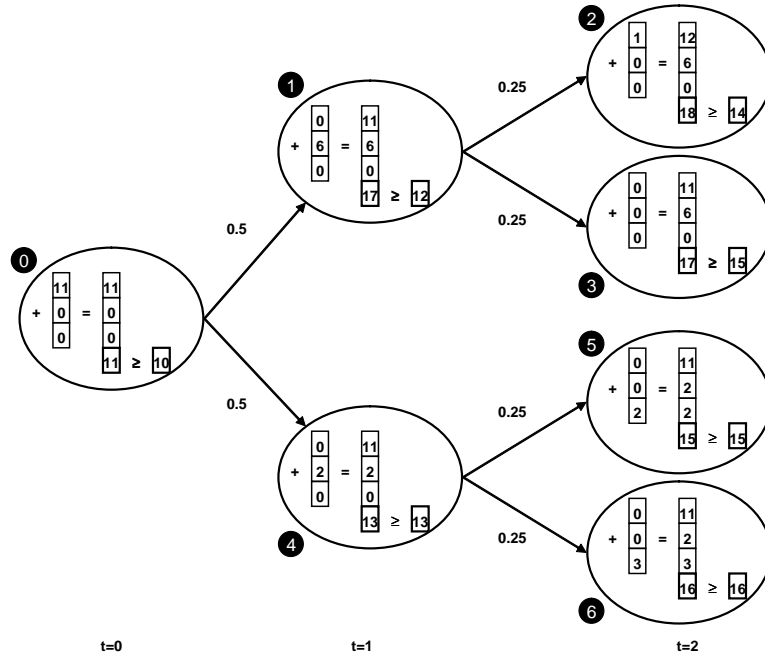


Figure 4.6: Solution before negative capacity variation

2. In node 2, the expansion is eliminated.
3. In node 1, no changes occur since there are no expansions and capacity is not lower than 10.
4. In node 3, no changes occur due to the fact that the reduction to 10 is compensated by the capacity surplus.
5. In node 0, the expansion is reduced from 11 to 10.
6. In node 4, the capacity expansions must be raised by 1. The resource where this expansion will have the lower cost is resource 2.

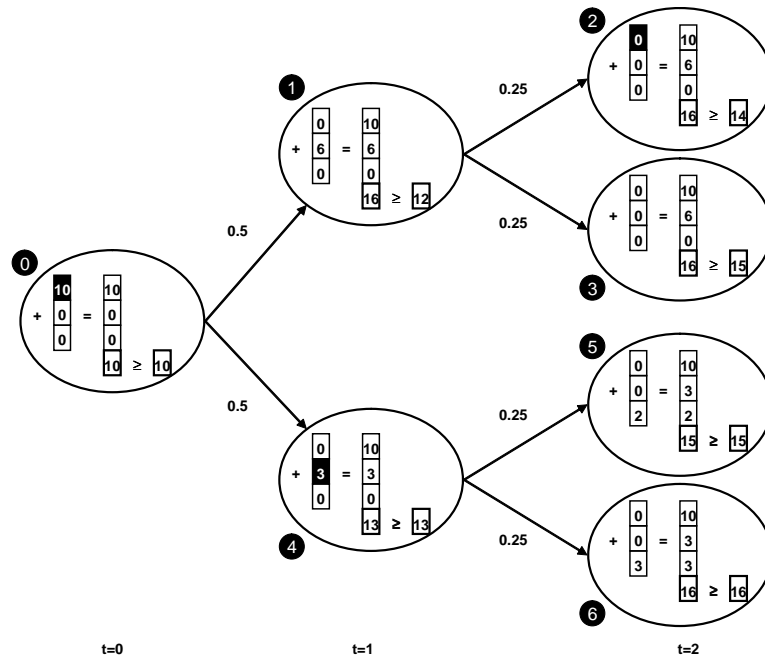


Figure 4.7: Solution after negative capacity variation

4.4.3 Expansion Shift (ES) Neighbourhood

This neighbourhood structure is defined by modifications to a solution consisting of partial or total shifts of capacity expansion from a selected resource in a selected node, to a resource in the same node or adjacent nodes:

- Backward - An expansion covering the reduction is created in the parent node for the resource that provides it for the minimum cost, and a procedure identical to the positive capacity variation is applied to the subtrees of the remaining children of the parent node.
- Same - An expansion covering the reduction and considering the node's capacity surplus is created in the same node for the other resource that provides it for the minimum cost.
- Forward - Reductions are limited by the demand satisfaction constraint for the selected node and a procedure identical to the negative capacity variation is used for the children of the selected node.
- Cut - Expansions are reduced by the selected node's capacity surplus.

4.4.4 Initial Solutions

Solutions are constructed by performing a depth first traversal of the tree, selecting a resource at each node and expanding its capacity by the minimum multiple of capacity increment that will enable satisfaction of the demand. The following criteria for *resource selection* and *definition of the demand to satisfy* have been defined:

- Resource selection can be made *randomly* or by a *cost minimisation* criterion.
- The *last tree level at which capacity expansions can occur* can be taken as a parameter, so that for the nodes above that level, expansions will cover the

node's demand, and for the nodes at that level, the maximum demand in the node's subtree will have to be satisfied - this allows for some flexibility in the definition of the timing of capacity expansion.

4.4.5 Objective Functions

Both expected value and CVaR_α have been implemented. A value of $\alpha = 0.9$ has been considered.

4.5 Computational Study

4.5.1 Test Instances

A set of problem instances was randomly generated for the purpose of this computational study. The guidelines for the creation of these instances are from the work by Ahmed et al. (2003) and by Huang and Ahmed (2005) (Table 4.1).

Table 4.1: Instance parameters

Parameter	Value
Branches for nonleaf nodes	3
Periods	4 (27 scenarios, 40 nodes) 5 (81 scenarios, 121 nodes)
Resources	3, 4
Capacity increments	Integer and uniform, between 1 and 5
Demand for root node, in $t = 0$ (d_0)	Integer and uniform, between 5 and 10
Fixed costs for root node, in $t = 0$ ($\alpha_{i,0}$)	Integer and uniform, between 2 and 10
Variable costs for root node, in $t = 0$ ($\beta_{i,0}$)	Integer and uniform, between 1 and 5
Demand and costs for a node in $t > 0$	Root value multiplied by lognormal variable Rounded to nearest integer Minimum of 1 for fixed and variable costs
Lognormal variable	4 patterns: $\delta_1 \sim \text{LN}(\mu = 1, \sigma = 0.5)$ $\delta_2 \sim \text{LN}(\mu = 1 + 0.5t, \sigma = 0.5)$ $\delta_3 \sim \text{LN}(\mu = 1, \sigma = 0.5 + 0.1t)$ $\delta_4 \sim \text{LN}(\mu = 1 + 0.5t, \sigma = 0.5 + 0.1t)$

One instance was created for each combination of number of periods, number of resources and evolution pattern for demand and costs, in a total of 16 instances. An ϵ -constraint method was used to obtain the nondominated sets for each instance. This method was implemented using the ILOG CPLEX 10.1 MIP solver. All experiments were performed in a platform with an Intel Xeon Dual Core 5160 3.0 GHz CPU, 8 GB RAM, running Red Hat Enterprise Linux 4. The software was generated with GCC 3.4.6 with level 3 optimisation.

Trees were implemented with *tree.hh*, an STL-like container class for n-ary trees, templated over the data stored at the nodes, developed by Kasper Peeters, and available at <http://www.damtp.cam.ac.uk/user/kp229/tree/>.

The nondominated sets were used to evaluate the quality of the approximations. For this evaluation we have chosen one of the unary quality indicators with fewer limitations: the hypervolume (Zitzler and Thiele, 1998) bounded by the set $(\mathbf{z}^1, \mathbf{z}^2, \dots)$ and a reference point (\mathbf{z}^{ref}) (Figure 4.8). For each instance, a reference point has been chosen so that all points in the nondominated and approximation sets lie in the hypervolume, by considering the worst values for each objective function degraded by an additional 0.1%. A relative measure was built upon this one, consisting of the ratio between the values of the indicators for the approximation set and the nondominated set, so as to enable comparison of performance across multiple instances. The quality gap indicator being used consists of the difference to 1 of this measure.

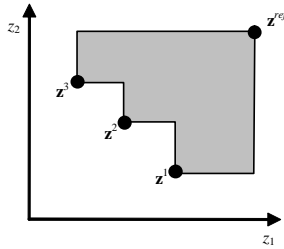


Figure 4.8: Hypervolume indicator

Considering, for example, the nondominated set $\{(1, 3), (2, 2), (3, 1)\}$ and an approximation set $\{(2, 4), (3, 3), (4, 2)\}$, for a problem with two minimisation objectives:

- the reference point is $(4.004, 4.004)$;
- the hypervolume of the nondominated set (shaded area in Figure 4.8) is 6.024016;
- the hypervolume of the approximation set is 1.016016;
- the quality gap indicator for the approximation set is $1 - 1.016016/6.024016 = 83.13\%$.

4.5.2 Algorithm Configuration

The computational experiments have been performed with an adaptation of TAMOCO, as implemented in MethOOD, with no tabu list and with fixed or variable sub-neighbourhoods. These configurations can be viewed as a Multiobjective Random Local Search, in the case of fixed sub-neighbourhoods, and a Multiobjective Variable Neighbourhood Search, in the case of variable sub-neighbourhoods.

With the results of a series of preliminary algorithm runs we have confined the range of parameter values to be studied to those presented in Table 4.2. The solution populations include 2 solutions constructed with each combination of the criteria for resource selection and last level for expansions.

Each of the 8 configurations has been run 30 times for each of the 16 instances. For all runs the generated approximation set has been recorded and its quality evaluated.

4.5.3 Experimental Results

In Table 4.3 we present the results obtained with the algorithmic configurations that provided higher quality results for the different neighbourhoods, individually - CV(50)

Table 4.2: Algorithm configurations

Parameter	Value
CV neighbourhoods	CV(50) - sub-neighbourhood of size 50% CV(75) - sub-neighbourhood of size 75%
ES neighbourhoods	ES(50) - sub-neighbourhood of size 50% ES(75) - sub-neighbourhood of size 75%
Variable neighbourhoods	CV(50)→ES(50) CV(50)→ES(75) CV(75)→ES(50) CV(75)→ES(75)
Population size	16 (problems with 4 periods); 20 (problems with 5 periods).
Time limit (seconds)	1 (problems with 4 periods); 5 (problems with 5 periods).

- and taken in combination - CV(75)→ES(50). In the identification of the problem instances we have *number of periods* / *number of resources* / *lognormal variable*.

Overall, the results are of high quality. Both configurations always find the efficient set in about half of the instances. For the remaining instances the efficient set is found in a large percentage of the runs while good approximations are obtained in the others. For each configuration, in only one instance does the average quality gap remain above 2%.

The use of variable neighbourhoods leads to significant improvement in the results in 4 instances and deterioration in 2 instances, with a globally positive effect in identical computational times. Computational times are at least an order of magnitude lower than the times required by the ϵ -constraint method, with the largest difference as high as 4 orders of magnitude.

4.6 Conclusions

The work described in this paper for multistage capacity expansion under uncertainty goes beyond what has been reported in the literature, by introducing an approach that considers lumpiness in capacity and both mean and risk criteria. MIP solvers can be used to obtain efficient sets for the problem, for example using the ϵ -constraint

Table 4.3: Computational study

Problem Instances	Quality Gap(%)				Time (seconds)				
	CV(50)		CV(75)→ES(50)		CV(50)		CV(75)→ES(50)		ϵ -constraint
	Mean	St. Dev.	Mean	St. Dev.	Mean	St. Dev.	Mean	St. Dev.	
4 / 3 / δ_1	0.00	0.00	0.00	0.00	0.01	0.00	0.01	0.00	0.81
4 / 3 / δ_2	0.00	0.00	0.00	0.00	0.00	0.00	0.00	0.00	0.55
4 / 3 / δ_3	0.00	0.00	0.00	0.00	0.05	0.04	0.07	0.09	3.97
4 / 3 / δ_4	0.00	0.00	0.02	0.07	0.14	0.11	0.19	0.23	3.06
4 / 4 / δ_1	0.00	0.00	0.00	0.00	0.02	0.02	0.01	0.01	5.88
4 / 4 / δ_2	0.00	0.00	0.00	0.00	0.01	0.01	0.01	0.01	24.93
4 / 4 / δ_3	0.04	0.20	0.57	0.58	0.27	0.18	0.26	0.19	2.41
4 / 4 / δ_4	1.79	3.93	0.27	0.52	0.11	0.13	0.17	0.24	21.97
5 / 3 / δ_1	0.00	0.00	0.00	0.00	0.07	0.02	0.08	0.04	7.28
5 / 3 / δ_2	0.02	0.00	0.02	0.01	0.31	0.12	0.30	0.19	3.91
5 / 3 / δ_3	0.88	0.93	0.46	0.66	2.49	1.47	2.20	1.13	459.05
5 / 3 / δ_4	0.07	0.37	2.24	4.25	2.19	1.42	0.99	1.27	19.75
5 / 4 / δ_1	0.00	0.00	0.00	0.00	0.16	0.11	0.09	0.04	588.67
5 / 4 / δ_2	2.08	2.30	0.04	0.20	1.51	1.01	1.76	1.06	92.56
5 / 4 / δ_3	0.00	0.00	0.00	0.00	0.07	0.01	0.09	0.02	1385.28
5 / 4 / δ_4	0.79	2.11	0.47	0.63	0.84	0.72	1.65	1.31	808.83

method. However, the computational times involved are very large, whereas the approach presented here can produce high quality approximations to the efficient sets in very short computational times.

Another unique feature of this approach is the fact that adopting different risk measures can be done by simply changing the corresponding objective function, while keeping the remaining parts of the implementation.

There are several ongoing extensions to this work: the consideration of flexible resources, forcing the inclusion of decisions on how to use the available resources to satisfy the demand; the exploitation of the properties of the problem to tackle the version with no lumpiness with metaheuristics; the study of features that solutions should exhibit in order to avoid stochastically dominated solutions in the efficient sets of mean-variance versions of the problem.

Chapter 5

A Multiobjective Metaheuristic for a Mean-Risk Multistage Capacity Investment Problem with Process Flexibility

(Under review at Journal of Scheduling)

In this paper, we propose a multiobjective local search metaheuristic for a mean-risk multistage capacity investment problem with process flexibility, irreversibility, lumpiness and economies of scale in capacity costs. In each period, discrete decisions concerning the investment in capacity expansion, and continuous decisions concerning the utilization of the available capacity to satisfy demand are considered. We solve the capacity utilization problems with linear programming, in order to find the minimum capacity for each resource with the other resources remaining unchanged, this way providing information on the feasibility of the discrete investment decisions.

Conditional value-at-risk is considered as a risk measure. Results of a computational study are presented, that show the approach is capable of obtaining high-quality approximations to the efficient sets, with a modest computational effort.

5.1 Introduction

In previous work (Claro and Sousa, 2007), we have addressed a problem of multi-stage capacity investment in dedicated resources with irreversibility, lumpiness and fixed-charge cost functions, in the presence of uncertainty. These features have been considered as main challenges in capacity modelling, in a recent survey on capacity management (Van Mieghem, 2003). Their relevance to industry applications has been emphasised by another recent survey on capacity management in the high-tech industry (Wu et al., 2005).

Here we extend that work by considering *process flexibility in the resources*, i.e., by allowing the possibility of producing different products in the same resource. This flexibility may be advantageous for firms that produce multiple products with uncertainty in the demand, if a favourable trade-off exists between the cost of dedicated resources and the benefits of demand-pooling (by taking advantage of changes or uncertainty in the relative proportions of the product quantities demanded) and a contribution margin option (by exploiting differentials in margin, i.e., the difference between price and variable cost, to produce and sell more of highly profitable products at the expense of less profitable products) (Van Mieghem, 1998). Previously, we proposed a multiobjective metaheuristic approach to tackle a bi-objective mean-risk formulation of the aforementioned problem, with conditional value-at-risk (CVaR) as a risk measure, and considering uncertainty in demand and capacity costs through a scenario tree. This problem is now extended to include *demand for multiple products* and *resources that may be used to produce multiple products*, and the metaheuristic-

tic approach is adapted to handle these features, in particular in what concerns the continuous decisions on the utilization of the available capacity.

The paper is organised as follows: section 5.2 reviews related work; section 5.3 presents a mixed integer programming (MIP) formulation for the problem, and a linear programming (LP) formulation for the capacity utilization problems; section 5.4 describes the multiobjective metaheuristic approach to the problem; section 5.5 presents the computational experiments and results; section 5.6 presents some conclusions and future work prospects.

5.2 Related Work

The starting point for our previous work (Claro and Sousa, 2007) was the problem of multistage dedicated capacity expansion under uncertainty considered in Ahmed et al. (2003). These authors use a scenario tree to model the stochastic evolution of costs and demand, and consider economies of scale in costs through fixed-charge cost functions. The survey on capacity management presented in Van Mieghem (2003) motivated us to readdress that problem, by including some additional challenging features in capacity modelling, namely lumpiness and the explicit consideration of risk. That work is now extended with the inclusion of *process flexibility in the resources*, in settings with *multiple product demands*.

Manufacturing flexibility has received significant attention in both literature and practice since the advent of flexible manufacturing systems in the 1980's. Reviews of research in manufacturing flexibility are available in Sethi and Sethi (1990), Kouvelis (1992), de Groote (1994), Toni and Tonchia (1998) and Beach et al. (2000).

Our work has partly been motivated by research in newsvendor models, in particular by Tomlin and Wang (2005) and Van Mieghem (2006), which consider risk-averse newsvendor networks. These two references provide reviews of the body of research

in both risk-averse single-resource newsvendors and newsvendor networks models, one of the major streams of literature on flexible capacity, essentially concerned with deriving managerial insights from those models.

A research stream that is closer to our work consists of mathematical programming approaches to investment in flexible capacity: Eppen et al. (1989) incorporate elements of scenario planning, integer programming and risk analysis in a model for strategic capacity planning in the automotive industry; Fine and Freund (1990) formulate and study a product-flexible capacity investment model as a two-stage nonlinear stochastic program; Li and Tirupati (1994) examine a multiproduct dynamic investment model for selection and expansion of dedicated and flexible capacity, formulated as a mathematical program, for which two heuristics are developed; Cheng et al. (2003) present a stochastic programming model for technology selection and capacity expansion with product mix flexibility and uncertainty in demand captured through a scenario-based approach, and develop a solution procedure based on an augmented Lagrangian method and restricted simplicial decomposition. Ahmed and Sahinidis (2003) formulate a stochastic capacity expansion problem with fixed-charge expansion cost functions and uncertainty in the problem parameters considered through scenarios as a multistage stochastic integer program, and develop a fast, linear-programming-based, approximation scheme.

Van Mieghem (2003) and Wu et al. (2005) present recent surveys on capacity management, and Julka et al. (2007) a review of capacity expansion, to which we direct interested readers. Claro and Sousa (2007) presents several references on the consideration of risk in stochastic optimisation and the application of (multiobjective) metaheuristics to stochastic optimisation problems and risk-return non-linear MIP problems. These are topics that are closely related to this work.

5.3 Mathematical Programming Models

We consider here a discretised planning horizon, over which the evolution of demands and costs is modelled with a scenario tree. Each level of the tree corresponds to a time period. $\mathcal{T}(n)$ denotes the subtree rooted in node n , with $n = 0$ being the root node, and $\mathcal{P}(n)$ the path from the root node to node n . \mathcal{S} is the set of leaf nodes, each corresponding to one of the S equally probable scenarios. In each node $n \in \mathcal{T}(0)$, each resource $i \in \mathcal{I}$ is characterised by a discounted fixed cost $\alpha_{i,n}$ and a discounted variable cost $\beta_{i,n}$ of expansion. The capacity of resource i can be changed by discrete increments of value l_i . The supply of capacity is unbounded and initial capacities are zero. The adaptations to include initial capacities are straightforward.

The demand for product $j \in \mathcal{J}$ is now given, in each node, by $d_{j,n}$. Each unit of capacity of resource i can produce $q_{i,j}$ units of product j . The decision variables are $x_{i,n}$, the number of capacity increments for resource i at node n , and $w_{i,j,n}$, the number of units of capacity of resource i allocated to the production of product j in node n . The binary variables $y_{i,n}$ take the value 1 if capacity of resource i is incremented in node n , and the value 0 otherwise.

Figures 5.1, 5.2 and 5.3 show a binary scenario tree and a feasible solution for a problem with 3 periods, 2 products ($j = 1, 2$), a resource dedicated to product 1 ($i = 1$), a resource dedicated to product 2 ($i = 2$), and 2 flexible resources ($i = 3, 4$):

- Figure 5.1 shows, at each node, capacity expansions and accumulated capacity, along with the capacity allocated to the production of each product in each resource. In each node, from left to right, the first column contains the capacity expansions performed for each resource in that node, the second column shows the capacity for each resource, the third column shows the amount of capacity allocated to the production of product 1, and the fourth to product 2.
- Figure 5.2 displays the solution's cost structure, with the two components con-

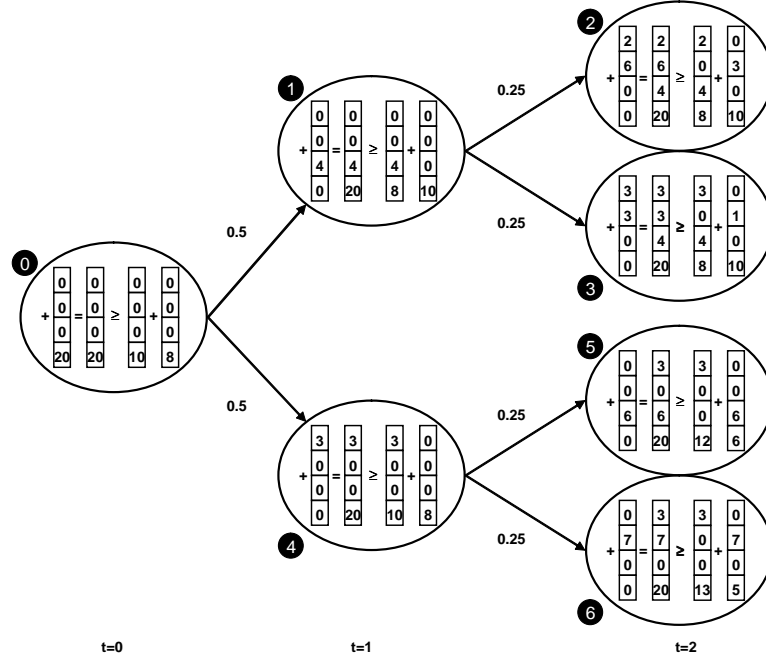


Figure 5.1: Solution and capacity constraints

sidered for each resource in each node - fixed and variable costs. In each node, from left to right, the columns contain: the fixed costs, the binary variables $y_{i,n}$, the variable costs and the capacity expansions.

- Figure 5.3 shows, for each node, how the total amount of each product that can be produced with the corresponding allocated capacity satisfies its demand. Each node is divided in two, horizontally: the upper part concerns product 1 and the lower product 2. In each part, the upper row contains the amount of capacity allocated in each resource to the production of the corresponding product, the lower row contains the amount of product each unit of allocated capacity can produce, the value on the left side of the inequality sign is the total amount of product, obtained by multiplying the allocated capacity by the amount of product each unit of capacity can produce, and the value on the right

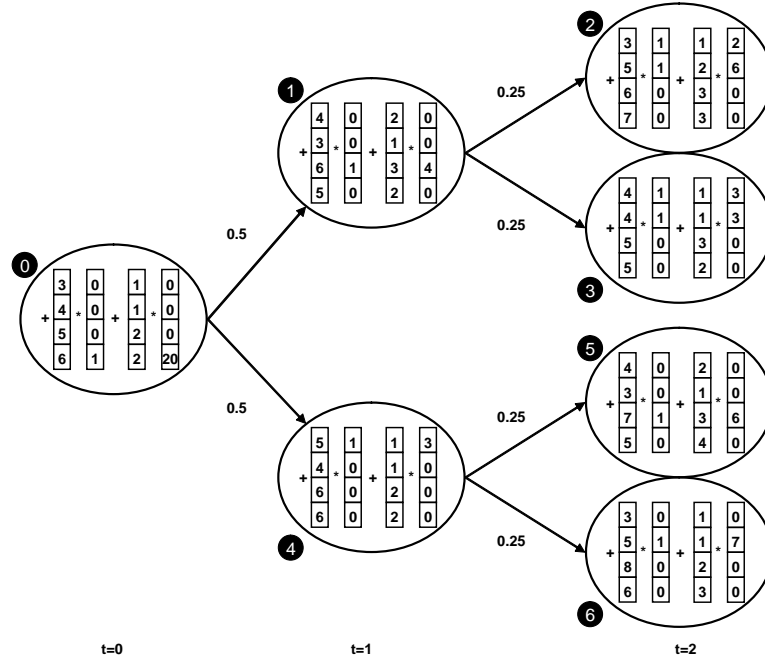


Figure 5.2: Capacity expansions and cost structure

side is the product demand in that node.

The solution is represented by the capacity expansions (Figures 5.1 and 5.2) and capacity utilization (Figures 5.1 and 5.3) in each node. To each arc is associated the probability of the node where the arc is directed to.

The conditional value-at-risk (CVaR_α), which can be viewed as the conditional expected value beyond value-at-risk (VaR_α), is a risk measure that has been receiving significant attention in the literature, mainly due to the fact that it is coherent and can be computed via linear programming (Rockafellar and Uryasev, 2002). Rigorous definitions of these measures can be found, for instance, in Rockafellar and Uryasev (2002). In a simplified way, we could say that, given a random variable representing cost, VaR_α could be defined as the maximum cost with probability level α , and CVaR_α as the expected value of the costs above that maximum cost with probability level α .

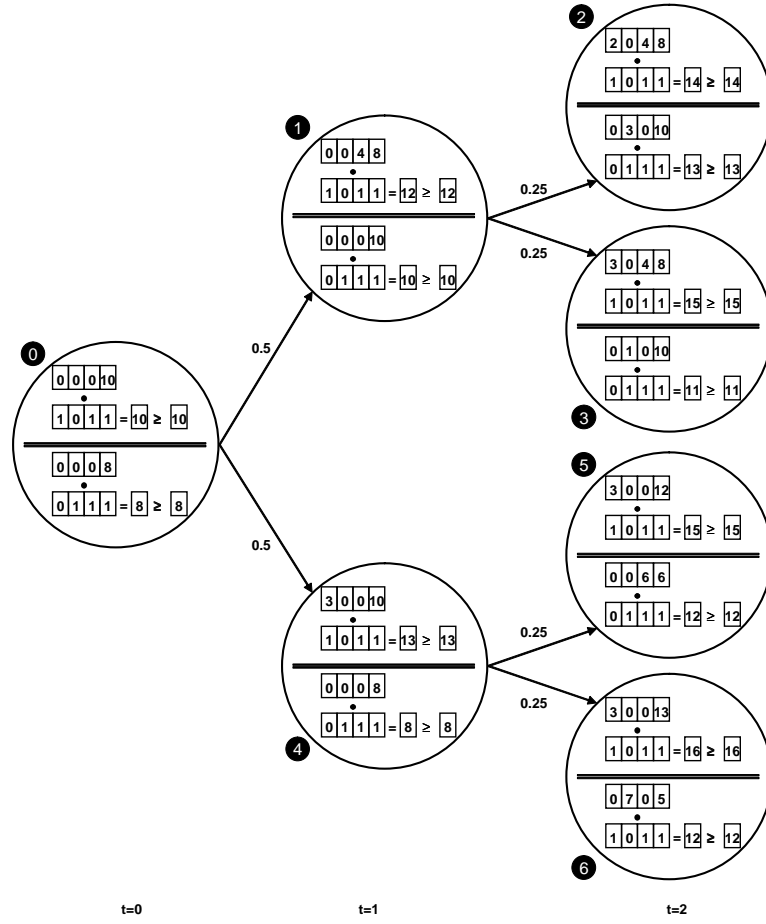


Figure 5.3: Capacity utilization and demand constraints

The consideration of a CVaR_α risk objective results in the following bi-objective mixed integer programming formulation:

$$\begin{aligned}
\min \quad & E = \frac{1}{S} \sum_{m \in \mathcal{S}} \sum_{n \in \mathcal{P}(m)} \sum_{i \in \mathcal{I}} (\alpha_{i,n} y_{i,n} + \beta_{i,n} l_i x_{i,n}) \\
\min \quad & \text{CVaR}_\alpha = \xi + \frac{1}{(1-\alpha)S} \sum_{m \in \mathcal{S}} Z_m \\
\text{s.t.} \quad & Z_m \geq \sum_{n \in \mathcal{P}(m)} \sum_{i \in \mathcal{I}} (\alpha_{i,n} y_{i,n} + \beta_{i,n} l_i x_{i,n}) - \xi, \quad m \in \mathcal{S} \\
& l_i x_{i,n} \leq M y_{i,n}, \quad n \in \mathcal{T}(0), i \in \mathcal{I} \\
& \sum_{i \in \mathcal{I}} q_{i,j} w_{i,j,n} \geq d_{j,n}, \quad n \in \mathcal{T}(0), j \in \mathcal{J} \\
& \sum_{j \in \mathcal{J}} w_{i,j,n} \leq \sum_{p \in \mathcal{P}(n)} l_i x_{i,p}, \quad n \in \mathcal{T}(0), i \in \mathcal{I} \\
& x_{i,n} \geq 0 \text{ and integer}, \quad n \in \mathcal{T}(0), i \in \mathcal{I} \\
& y_{i,n} \in \{0, 1\}, \quad n \in \mathcal{T}(0), i \in \mathcal{I} \\
& w_{i,j,n} \geq 0, \quad n \in \mathcal{T}(0), i \in \mathcal{I}, j \in \mathcal{J} \\
& Z_m \geq 0, \quad m \in \mathcal{S} \\
& \xi \geq 0.
\end{aligned} \tag{5.1}$$

In this formulation, the objective E minimises the *expected value of the total investment cost* over the planning horizon, and the objective CVaR_α minimises the *conditional value-at-risk of that investment cost*. The first set of constraints are the usual constraints required for defining the conditional value-at-risk. The second set of constraints define $y_{i,n}$ in terms of variables $x_{i,n}$. The third and fourth sets of constraints are the demand satisfaction and capacity constraints, respectively.

A useful information regarding the feasibility of changes to the discrete investment decisions in a given solution, is the minimum capacity $K_{r,n}^*$ that is required for a resource r in a node n , considering only the conditions of node n and with the capacities of all other resources remaining unchanged. To compute this minimum capacity $K_{r,n}^*$ we define a new problem, where the variables are $K_{r,n}$ and $w_{i,j,n}$, $i \in$

$\mathcal{I}, j \in \mathcal{J}$, while $x_{i,p}, p \in \mathcal{P}(n), i \in \mathcal{I} \setminus \{r\}$ become problem parameters.

The new problem can be formulated as a linear program, as follows:

$$\begin{aligned}
& \min && K_{r,n} \\
& \text{s.t.} && \sum_{i \in \mathcal{I}} q_{i,j} w_{i,j,n} \geq d_{j,n}, && j \in \mathcal{J} \\
& && \sum_{j \in \mathcal{J}} w_{i,j,n} \leq \sum_{p \in \mathcal{P}(n)} l_i x_{i,p}, && i \in \mathcal{I} \setminus \{r\} \\
& && \sum_{j \in \mathcal{J}} w_{r,j,n} \leq K_{r,n}, \\
& && w_{i,j,n} \geq 0, && i \in \mathcal{I}, j \in \mathcal{J} \\
& && K_{r,n} \geq 0.
\end{aligned} \tag{5.2}$$

The first set of constraints in this formulation are the demand satisfaction constraints. The other constraints model capacities.

5.4 A Multiobjective Metaheuristic Approach

Multiobjective metaheuristics have been successfully applied to Multiobjective Combinatorial Optimisation (MOCO) problems and are particularly well-suited to deal with the set of previously mentioned challenging features in capacity modelling. Surveys on multiobjective metaheuristics are available in Ehrgott and Gandibleux (2000) and Jones et al. (2002).

Tabu Search for Multiobjective Combinatorial Optimisation (TAMOCO) (Hansen, 2000) and Pareto Simulated Annealing (PSA) (Czyzak and Jaskiewicz, 1998) can be viewed as Multiobjective Local Search (MOLS) approaches. Both aim at producing a good approximation of the efficient set, working with a population of solutions, each solution holding a weight vector for the definition of a search direction. Each approach proposes a different strategy for the definition of the weights, but with

the same purpose: orientation of the search towards the nondominated frontier and spreading of solutions over that frontier (the former is achieved by the use of positive weights, while the latter is based on a comparison with other solutions of the population). Although in different ways, both methods operate on each single solution, searching and selecting a solution in its neighbourhood that will eventually replace it. Moreover, each procedure involves traditional metaheuristic components such as neighbourhoods, in general, or tabu lists, in the specific case of TAMOCO. The identification of these common aspects has suggested the definition of a MOLS generic template (Algorithm 7).

PSA and TAMOCO differ in the definition of several of the template's primitive operations: weight vectors are distinctly initialised and updated; $Neighbourhood(s)$ in PSA is a random subneighbourhood with just one movement; the generated movement in PSA is always selected, while in TAMOCO movement selection considers tabu status, aspiration criteria, and a comparison of evaluations based on a weighted sum scalarising function; in PSA a selected movement is accepted according to an

Algorithm 7: Multiobjective Local Search Template

```

Generate a set of initial feasible solutions  $G \subset S$ ;
Initialise the approximation to the efficient set  $E = \{\}$ ;
foreach  $s_i \in G$  do
    Initialise the corresponding context;
    Update  $E$  with  $s_i$ ;
end
while a stopping criterion is not met do
    foreach  $s_i \in G$  do
        Update the corresponding weight vector  $\lambda_i$ ;
        Initialise the selected solution  $s_s = 0$ ;
        foreach  $s' \in Neighbourhood(s_i)$  do
            Update  $E$  with  $s'$ ;
            if  $s'$  is selectable and  $s'$  is preferable to  $s_s$  then  $s_s = s'$ ;
        end
        if  $s_s \neq 0$  and  $s_s$  is acceptable then  $s_i = s_s$ ;
    end
end

```

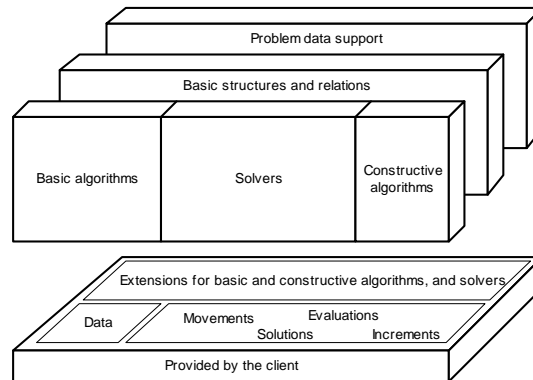


Figure 5.4: The MetHOOD framework

acceptance probability, while in TAMOCO it is always accepted.

This template and related procedures have been implemented in an object-oriented framework called MetHOOD (Claro and Sousa, 2001) (Figure 5.4), that has been used to support the application described in this paper. Also of interest for this application is the support that the framework provides for neighbourhood variation, i.e., we can consider a sequence of neighbourhood structures and use them dynamically according to the evolution of the search process: if a new accepted solution is preferable to the current one, or if the current neighbourhood is the last in the sequence, the first neighbourhood in the sequence will be used next; otherwise the following neighbourhood in the sequence will be used next.

The MetHOOD framework has been instantiated for the capacity investment problem according to the implementation choices described next. With this framework instantiation several MOLS algorithms become readily available.

5.4.1 Basic Concepts

The instantiation of the MetHOOD framework involved the definition of several concepts derived from the basic description of the problem:

- A resource's *capacity equivalent* is a set of resources that is able to produce the resource's range of products and has no *redundant* resources, i.e., resources whose exclusion does not eliminate the ability to produce any product in that range. For the instance considered in Figures 5.1, 5.2 and 5.3, the equivalents are the following: for resource 1, $\{1\}$, $\{3\}$ and $\{4\}$; for resource 2, $\{2\}$, $\{3\}$ and $\{4\}$; for resource 3, $\{1, 2\}$, $\{3\}$ and $\{4\}$; for resource 4, $\{1, 2\}$, $\{3\}$ and $\{4\}$.
- A *capacity out-shift factor* is defined for any ordered pair of resources (i_1, i_2) that share at least one product, and consists of the maximum $q_{i_1,j}/q_{i_2,j}$ ratio over all shared products. Multiplying the capacity of i_1 by the out-shift factor results in just enough capacity of i_2 to produce the maximum quantity that i_1 would be able to produce of any of the shared products. For the instance considered in Figures 5.1, 5.2 and 5.3, all factors are 1.
- In each node n , an array with the solutions to problem (5.2) for each resource i , $K_{i,n}^*$ may be defined, as well as an array containing the excesses of capacity over each $K_{i,n}^*$. Based on these, arrays of *forward capacity surplus* (the minimum excess in the node's subtree) and *backward capacity surplus* (the capacity in excess from the maximum $K_{i,p}^*$ in the nodes belonging to the path from the root to the node) may also be defined. For the problem instance and solution described in Figures 5.1, 5.2 and 5.3, Figure 5.5 displays, for each node, from left to right, arrays of capacity, $K_{i,n}^*$, excess of capacity over $K_{i,n}^*$, backward and forward capacity surplus.

In node 1, for example, only flexible resources 3 and 4 have capacity: 4 units for resource 3, and 20 for resource 4. The demands are 12 for product 1, and 10 for product 2. For the flexible resource 4:

- Since $q_{3,1} = q_{3,2} = q_{4,1} = q_{4,2} = 1$, to compute $K_{4,1}^*$ it is indifferent which products are assigned to which resources. From the total demand of 22, a

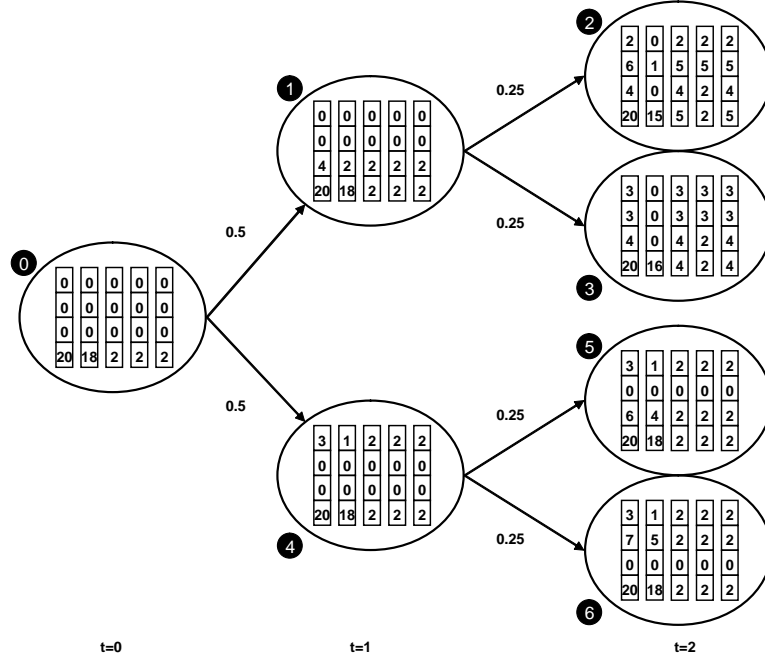


Figure 5.5: Capacity, $K_{i,n}^*$, excess of capacity over $K_{i,n}^*$, backward and forward capacity surplus

maximum of 4 can be assigned to resource 3, and the remaining 18 have to be assigned to resource 4. $K_{4,1}^*$ is, therefore, 18.

- The excess of capacity is $20 - 18 = 2$.
- The forward capacity surplus is the minimum excess, for resource 4, in nodes 1, 2 and 3, i.e., $\min \{2, 5, 4\} = 2$.
- The backward capacity surplus is the capacity in excess of the maximum $K_{4,n}^*$, in nodes 0 and 1, i.e., $20 - \max \{18, 18\} = 2$.

5.4.2 Solution

A tree representation has been used, matching the scenario tree structure that models demand and costs (see Section 5.3), and considering arrays of capacity expansions and

LP models (5.2) for each resource in each node. Subtrees entirely consisting of nodes with a demand vector lower than or equal to at least one demand vector in the path from root to their parent are not considered in the solution representation, since no further expansions are required to satisfy the demand in those nodes.

5.4.3 Capacity Variation (CV) Neighbourhood

This neighbourhood structure is defined by modifications to a solution consisting of positive or negative capacity variations, for a selected resource in a selected node.

Positive Variations

We consider as an upper limit on positive variations the capacity expansion that would be required for the resource to satisfy the joint demand of all the products it can produce, in the node and its subtree, reduced by the demand that can be satisfied by the dedicated resources in the parent node.

Considering the positive variation, the node (only for the remaining resources) and the other nodes in its subtree, are visited in depth-first order and the maximum feasible reductions (the forward capacity surpluses) of capacity expansions are performed. If a node has more than one possible expansion, the resource with higher cost reduction is chosen first. After each reduction, the forward capacity surpluses are recomputed.

For the problem instance and solution described in Figures 5.1, 5.2, 5.3 and 5.5, Figure 5.6 illustrates a positive capacity variation for resource 4 in node 0, from 20 to 22. In the figure, which has a structure identical to Figure 5.1, the boxes with black background correspond to the adjustments in capacity expansions:

1. In node 0, there are no other capacity expansions.
2. In node 1, due to the capacity variation in node 0, the forward capacity surplus

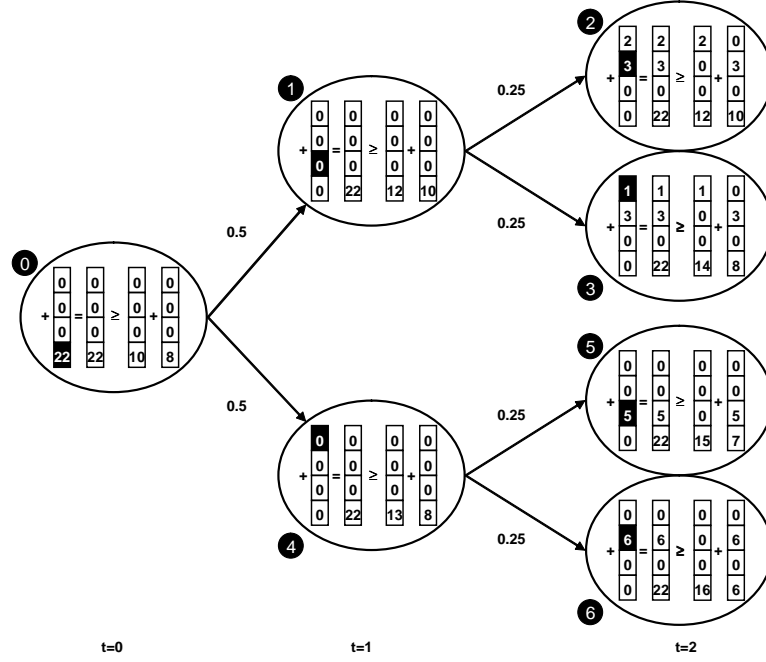


Figure 5.6: Solution after positive capacity variation

for resource 3 changes to 4. The capacity expansion has a value of 4, not higher than the capacity surplus, so it can be discarded.

3. In node 2, due to the capacity variations in nodes 0 and 1, the forward capacity surpluses for resources 1 and 2 change to 2 and 3, respectively. The higher cost reduction is achieved for resource 2, so its capacity expansion is reduced by 3, to 3.
4. In node 3, due to the capacity variations in nodes 0 and 1, the forward capacity surpluses for both resources 1 and 2 change to 2. The cost reductions are identical, so the capacity expansion in resource 1 is reduced by 2, to 1.
5. In node 4, due to the capacity variation in node 0, the forward capacity surplus for resource 1 changes to 3, so the entire expansion is discarded.

6. In node 5, due to the capacity variations in nodes 0 and 4, the forward capacity surplus for resource 3 changes to 1, so its capacity expansion is reduced by 1, to 5.
7. In node 6, due to the capacity variations in nodes 0 and 4, the forward capacity surplus for resource 2 changes to 1, so its capacity expansion is reduced by 1, to 6.

Negative Variations

Negative variations are limited by the backward capacity surpluses. The capacity level has to be adjusted for the set of nodes in the path from the selected node to the root where the selected resource's capacity level is higher or equal to the reduced value: in the topmost of these nodes, the resource's expansion is reduced so that its capacity becomes equal to the reduced value; in the remaining nodes, existing expansions of the selected resource are discarded.

In the children of the nodes in this set that are not themselves part of the set, expansions covering the capacity reduction in the parent node are created for the capacity equivalent that provides them at the minimum cost. The set of expansions for a capacity equivalent is found by, first, expanding the capacity of all its resources by the minimum amount that covers the capacity reduction (obtained with the application of the relevant out-shift factor), and then, iteratively reducing the expansions as much as possible starting with the ones that present higher cost reductions.

For the problem instance and solution described in Figures 5.1, 5.2, 5.3 and 5.5, Figure 5.7 illustrates a negative capacity variation for resource 4 in node 2, from 20 to 18. In the figure, which has a structure identical to Figure 5.1, the boxes with black background correspond to the adjustments in capacity expansions:

1. The capacity reduction for resource 4 in node 2 is limited by its backward

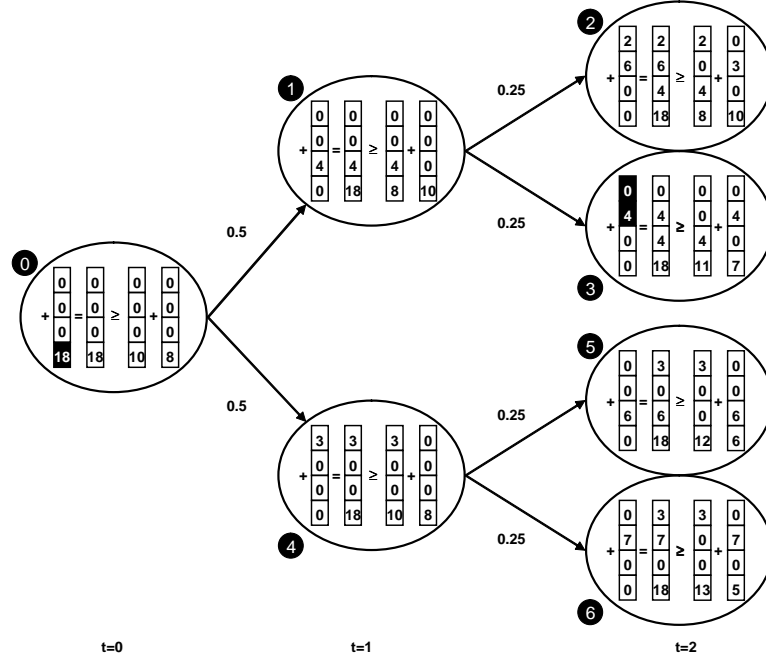


Figure 5.7: Solution after negative capacity variation

capacity surplus, with a value of 2. The capacity can not, therefore, be reduced below 18. The proposed reduction is within this range.

2. In node 1, no changes occur since there are no expansions and capacity is not lower than 18.
3. In node 3, an adjustment of capacity may be required, due to the capacity reduction in its parent node. The candidates for this adjustment are the capacity equivalents of resource 4:

- For $\{1, 2\}$, the expansions in both resources are increased by 2, and could then be reduced by 5. Both reductions have the same cost, so the expansion in resource 1 is discarded and the expansion in resource 2 is increased by 1, to 4.

- For $\{3\}$, the expansion is increased by 2, and could be reduced by 2.
- For $\{4\}$, the expansion is increased by 2, and could be reduced by 2.

The changes in capacity equivalent $\{1, 2\}$ have the lower cost (negative, in this case), so they are performed.

4. In node 0, the capacity expansion of resource 4 is reduced from 20 to 18.
5. In node 4, an adjustment of capacity may be required, due to the capacity reduction in its parent node. The candidates for this adjustment are, again, the capacity equivalents of resource 4:

- For $\{1, 2\}$, the expansions in resources 1 and 2 are both increased by 2, and could then be reduced by 4 and 2, respectively. The achievable cost reduction is higher for resource 2, so the expansion in resource 2 remains 0, and the expansion in resource 1 remains 3.
- For $\{3\}$, the capacity is increased by 2, and could be reduced by 2.
- For $\{4\}$, the capacity is increased by 2, and could be reduced by 2.

All alternatives lead to no changes in the expansions in this node.

5.4.4 Expansion Shift (ES) Neighbourhood

This neighbourhood structure is defined by modifications to a solution consisting of capacity expansion cuts and partial or total shifts of capacity expansions in the same node:

- Cut - Expansions are reduced by the selected resource's forward capacity surplus.

- Same - As in negative variations, expansions covering the new capacity requirements are created for the capacity equivalent that provides them at the minimum cost, with the exception of the origin resource. The set of expansions for a capacity equivalent is found with the same logic as in negative variations.

5.4.5 Initial Solutions

Solutions are constructed by performing a depth first traversal of the tree. At each node, a sequence of capacity expansions are performed for a random sequence of the products. This sequence of expansions is computed with the use of the LP models (5.2) for the resources in the node:

- Initial conditions for all models are *no demand* and *disabling of all capacity utilization*.
- At each iteration of the sequence of products:
 - The demand for that product is introduced in the LP models, and capacity utilization for that product is enabled.
 - For each resource that is able to produce the product, the LP model is solved in order to determine its minimum required expansion. The expansion with lower cost is performed.

The last tree level at which capacity expansions can occur can be taken as a parameter, so that for the nodes above that level, expansions will cover the node's demand, and for the nodes at that level, the maximum demands for each product in the node's subtree will have to be satisfied. This allows for some flexibility in the definition of the timing of capacity expansion.

5.4.6 Objective Functions

Both expected value and CVaR_α have been implemented. A value of $\alpha = 0.9$ has been considered.

5.5 Computational Study

5.5.1 Instances

A set of problem instances was generated for the purpose of this computational study. The guidelines for the creation of these instances are from the work described by Ahmed et al. (2003), Huang and Ahmed (2005) and Claro and Sousa (2007) (Table 5.1). Straightforward adaptations were made to include the demands for 2 products, that were considered to be independent, and to define the amount of product produced by unit of capacity, for which values of either 0 or 1 were considered. In the root node, the variable costs for flexible resources are forced to be higher than the variable costs for dedicated resources and lower than the sums of variable costs for dedicated resources. A total of 10 instances were created, 5 for each number of periods.

For the instances with 4 time periods, an ϵ -constraint method was used to obtain the nondominated sets. For the instances with 5 time periods, the single objective constrained problems that are solved in the ϵ -constraint method became very hard to solve, requiring computational times in the order of thousands of seconds just to find a feasible solution. Our option was, in this case, to use weighted sum formulations of the bi-objective problem to obtain reasonable approximations to the efficient sets. These formulations are computationally less demanding; however, they are unable to find nonsupported efficient solutions and even some supported efficient solutions may be missed if the intervals between the weights considered in each problem are not

Table 5.1: Instance parameters

Parameter	Value
Branches for nonleaf nodes	3
Periods	4 (27 scenarios, 40 nodes) 5 (81 scenarios, 121 nodes)
Products	2
Resources	2 flexible and 2 dedicated (1 for each product)
Capacity increments	Integer and uniform, between 1 and 5
Capacity utilization	1 for all feasible product-resource pairs
Demands for root node, in $t = 0$ (d_0)	Integer and uniform, between 5 and 10
	Independent
Fixed costs for root node, in $t = 0$ ($\alpha_{i,0}$)	Integer and uniform, between 2 and 10
Variable costs for root node, in $t = 0$ ($\beta_{i,0}$)	Integer and uniform, between 1 and 5
Demand and costs for a node in $t > 0$	Root value multiplied by lognormal variable $\delta \sim \text{LN}(\mu = 1 + 0, 5t, \sigma = 0.5 + 0.1t)$ Rounded to nearest integer Minimum of 1 for fixed and variable costs

tight enough. Range equalisation was not used because the expected value and the CVaR_α have values in the same order of magnitude. A step of 0.05 was used for the weights and a time limit of 1000 seconds was imposed for each problem. The ILOG CPLEX 10.1 MIP solver was used.

All experiments were performed in a platform with an Intel Xeon Dual Core 5160 3.0 GHz CPU, 8 GB RAM, running Red Hat Enterprise Linux 4. The software was generated with GCC 3.4.6 with level 3 optimisation.

Trees were implemented with *tree.hh*, an STL-like container class for n-ary trees, templated over the data stored at the nodes, developed by Kasper Peeters, and available at <http://www.damtp.cam.ac.uk/user/kp229/tree/>. The LP problems are solved with the ILOG CPLEX 10.1 LP solver.

The reference sets obtained with the MIP solver were used to evaluate the quality of the approximations. For this evaluation we have chosen one of the unary quality indicators with fewer limitations: the hypervolume (Zitzler and Thiele, 1998) bounded by the set $(\mathbf{z}^1, \mathbf{z}^2, \dots)$ and a reference point (\mathbf{z}^{ref}) (Figure 5.8). For each instance, a reference point has been chosen so that all points in the reference and approxima-

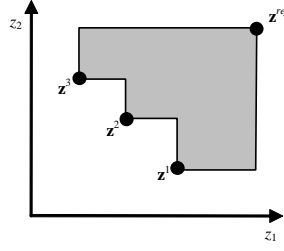


Figure 5.8: Hypervolume indicator

tion sets lie in the hypervolume, by considering the worst values for each objective function degraded by an additional 0.1%. A relative measure was built upon this one, consisting of the ratio between the values of the indicators for the approximation set and the reference set, so as to enable comparison of performance across multiple instances. The quality gap indicator being used consists of the difference to 1 of this measure.

Considering, for example, the reference set $\{(1, 3), (2, 2), (3, 1)\}$ and an approximation set $\{(2, 4), (3, 3), (4, 2)\}$, for a problem with two minimisation objectives:

- the reference point is $(4.004, 4.004)$;
- the hypervolume of the reference set (shaded area in Figure 5.8) is 6.024016;
- the hypervolume of the approximation set is 1.016016;
- the quality gap indicator for the approximation set is $1 - 1.016016/6.024016 = 83.13\%$.

5.5.2 Algorithm Configuration

These computational experiments have been performed with an adaptation of TAMOCO, as implemented in MetHOOD, with no tabu list and with fixed or variable sub-neighbourhoods. These configurations can be viewed as a Multiobjective Random

Table 5.2: Algorithm configurations

Parameter	Value
CV neighbourhood	CV(15) - sub-neighbourhood of size 15%
ES neighbourhood	ES(15) - sub-neighbourhood of size 15%
Variable neighbourhood	CV(15)→ES(15)
Population size	10 (problems with 4 periods); 15 (problems with 5 periods).
Time limit (seconds)	120 and 300 (problems with 4 periods); 1800 and 3600 (problems with 5 periods).

Local Search, in the case of fixed sub-neighbourhoods, and a Multiobjective Variable Neighbourhood Search, in the case of variable sub-neighbourhoods.

With the results of a series of preliminary algorithm runs we have confined the range of parameter values to be studied to those presented in Table 5.2. The solution populations include an increasing number of solutions constructed with increasing parameter values for the last level of expansion: 1 for level 0, 2 for level 1, 3 for level 2, 4 for level 3, and 5 for level 4 (for the problems with 5 periods).

Each configuration has been run 30 times for each instance. For all runs the generated approximation set has been recorded and its quality evaluated.

5.5.3 Experimental Results

The quality gap indicator and the computational times for the experimental results are presented in Tables 5.3, 5.4, 5.5 and 5.6.

The configuration with best overall performance is CV(15)→ES(15) at the longest running times: for the instances with 4 periods, the average quality gap is always below 1%, with average computational times that range from equivalent to largely less than those used by the ϵ -constraint method; for the instances with 5 periods, the average approximation quality gap is consistently below the results obtained with the weighted sum approach, with average computational times that are less than half the computational times used for the weighted sum approach.

Table 5.3: Computational study for instances with 4 periods and time limit 300 s

Instance	Quality Gap(%)				Time (seconds)				
	CV(15)		CV(15)→ES(15)		CV(15)		CV(15)→ES(15)		ϵ -constraint
	Mean	St. Dev.	Mean	St. Dev.	Mean	St. Dev.	Mean	St. Dev.	
4_1	0.40	0.39	0.38	0.40	236.91	61.37	212.43	58.30	642.81
4_2	0.56	0.55	0.65	0.56	199.42	81.39	195.30	76.67	675.69
4_3	0.00	0.00	0.00	0.00	74.88	45.99	34.38	22.34	59329.68
4_4	1.04	0.00	0.98	0.18	59.48	35.82	88.07	63.07	1773.56
4_5	0.18	0.14	0.15	0.13	101.04	78.88	133.82	83.99	129.28

Table 5.4: Computational study for instances with 4 periods and neighbourhood CV(15)→ES(15)

Instance	Quality Gap(%)				Time (seconds)				
	120 s		300 s		120 s		300 s		ϵ -constraint
	Mean	St. Dev.	Mean	St. Dev.	Mean	St. Dev.	Mean	St. Dev.	
4_1	0.98	0.53	0.38	0.40	95.48	20.50	212.43	58.30	642.81
4_2	1.05	0.63	0.65	0.56	72.29	26.33	195.30	76.67	675.69
4_3	0.00	0.00	0.00	0.00	41.22	26.04	34.38	22.34	59329.68
4_4	1.08	0.14	0.98	0.18	61.37	31.30	88.07	63.07	1773.56
4_5	0.28	0.10	0.15	0.13	52.63	37.16	133.82	83.99	129.28

Two statistical comparisons between algorithm configurations were performed: for the variable neighbourhood configuration, between shorter and longer running times; for the longer running times, between single and variable neighbourhood configurations. These comparisons were performed separately for the group of instances with 4 time periods and the group of instances with 5 time periods. The data violates ANOVA assumptions, so pairwise comparisons of the algorithms were performed for each instance, using Mann-Whitney tests, and the individual significance levels were adjusted using a Bonferroni correction to provide an overall significance level of 0.05. For each group of 5 instances, 10 comparisons were performed, resulting in an individual significance level of 0.005.

Comparing the single and variable neighbourhood configurations, with longer running times, statistically significant differences in approximation quality were found for only one instance with 5 periods. No other significant differences were found.

In the group of instances with 4 periods, the increase in available running time resulted in statistically significant improvements for 3 instances. For the remaining

Table 5.5: Computational study for instances with 5 periods and time limit 3600 s

Instance	Quality Gap(%)				Time (seconds)				
	CV(15)		CV(15)→ES(15)		CV(15)		CV(15)→ES(15)		Weighted Sum
	Mean	St. Dev.	Mean	St. Dev.	Mean	St. Dev.	Mean	St. Dev.	
5_1	-4.17	0.52	-4.34	0.45	3182.35	446.22	3302.17	366.65	6987.27
5_2	0.19	0.70	-0.97	0.59	3512.45	158.89	3479.32	171.02	13424.90
5_3	-0.74	1.85	-1.19	1.47	2953.42	644.63	3111.35	492.69	6635.44
5_4	-11.53	3.96	-13.38	2.66	2694.72	771.68	2905.17	722.74	11190.90
5_5	-3.94	1.93	-4.69	1.26	2786.22	669.36	3082.48	547.46	9523.72

Table 5.6: Computational study for instances with 5 periods and neighbourhood CV(15)→ES(15)

Instance	Quality Gap(%)				Time (seconds)				
	1800 s		3600 s		1800 s		3600 s		Weighted Sum
	Mean	St. Dev.	Mean	St. Dev.	Mean	St. Dev.	Mean	St. Dev.	
5_1	-3.57	0.85	-4.34	0.45	1730.49	86.48	3302.17	366.65	6987.27
5_2	0.27	1.14	-0.97	0.59	1792.11	34.31	3479.32	171.02	13424.90
5_3	0.15	1.55	-1.19	1.47	1669.71	177.87	3111.35	492.69	6635.44
5_4	-12.51	3.11	-13.38	2.66	1483.94	268.58	2905.17	722.74	11190.90
5_5	-4.08	1.56	-4.69	1.26	1709.30	193.76	3082.48	547.46	9523.72

instances (4_3 and 4_4) an early convergence, as can be observed by the similarity in convergence times, justifies the absence of differences. In the group of instances with 5 periods, the increase in available computational time resulted in statistically significant improvements also for 3 instances. For the remaining instances (5_4 and 5_5) a large dispersion of the results renders the quality improvement statistically insignificant.

5.6 Conclusions

Our previous work (Claro and Sousa, 2007) went beyond work reported in the literature regarding multistage capacity expansion under uncertainty, by presenting an approach that considered lumpiness in capacity and both mean and risk criteria. Here we have extended that work here by considering process flexibility in the resources. The resulting problem is a MIP problem that we have addressed with a multiobjective local search metaheuristic, integrating linear programming to tackle the continuous

decisions on the utilization of the available flexible capacity to satisfy demand.

MIP solvers can be used to obtain efficient sets for this problem, for example using the ϵ -constraint method. The computational times, however, soon become prohibitive, whereas the approach presented here leads to high quality approximations to the efficient sets in comparatively reduced computational times.

The incorporation of flexible capacity is a first step to enrich this capacity expansion model. In future work we will be aiming at incorporating product prices, processing costs and capacity leadtimes in the model. Another line of development that we will be looking into is the integration of this optimisation model with a game-theoretic framework to tackle capacity expansion in multiresource multiagent networks.

Chapter 6

Conclusion

6.1 Background

Many important decisions in Operations Management, in particular at a strategic level, are made in the presence of uncertainty, and should naturally consider the variability of their results. However, the majority of research and applied work in this area ignores this aspect and focuses exclusively on an expected value decision criterion. These decision settings are usually complex and if mathematical models are to be used to characterise them, they are required to include uncertainty in the parameters, logical and other discrete decision variables, and multiple objectives.

One of the critical decision areas within Operations Strategy is Capacity Expansion, which is concerned with deciding the type, magnitude, timing, and location of capacity acquisition. Capacity models are required to address a variety of problem features related to the previously mentioned complexity, leading to nonlinearities, nonconvexities, integrality and multiple objectives.

Multiobjective metaheuristics are optimisation algorithms extremely well suited to efficiently tackle problems that present these difficulties and have therefore the potential to play an important role as general approaches for mean-risk combina-

torial optimisation problems. The primary objective of our work was to perform a preliminary assessment of this potential.

6.2 Main Contributions

Static Stochastic Knapsack

We have proposed an approach that:

- is able to handle both exact and sample approximation formulations of the problem, and extends these formulations to consider both mean and risk criteria;
- produces high quality approximations to the efficient sets in computational times much shorter than state-of-the-art IP and QIP/QCIP solvers;
- can easily accommodate different risk measures by simply changing the implementation of the corresponding objective function.

Multistage Capacity Investment

The approach developed for this problem:

- extends previous formulations proposed in the literature, by considering lumpiness in capacity expansions and both mean and risk criteria;
- produces high quality approximations to the efficient sets in computational times much shorter than state-of-the-art IP solvers.

Multistage Capacity Investment with Process Flexibility

We have developed an approach for this problem that:

- extends our previous formulation by considering process flexibility in the resources;
- addresses the continuous decisions in a MIP problem integrating linear programming in a multiobjective local search metaheuristic;
- can obtain high quality approximations to the efficient sets in comparatively reduced computational times.

General Contributions

In a more general perspective, we can view the main contributions of the work presented in this dissertation as being the following:

- we explicitly introduce a multiobjective mean-risk framework for the general class of Stochastic Combinatorial Optimisation problems;
- we propose and assess the potential of multiobjective metaheuristics as a class of algorithms well suited to deal with the difficulties presented by these problems.

6.3 Future Developments

In the broader scope of decision making in Operations Strategy, Capacity Expansion problems are probably the application area where our work may have a larger impact.

In a near future, the developments of this work should focus on:

- the enrichment of the capacity expansion model, incorporating product prices, processing costs and capacity leadtimes;
- integrating the multiobjective optimisation framework with a game-theoretic framework to tackle multiresource multiagent networks, this being a natural

extension of capacity investment models since capacity decisions often are contingent upon the decisions of the other economic agents (customers, suppliers, partners, competitors) in the network.

Alongside these developments, two additional research topics, of a more practical nature, should be considered:

- one is the task of deriving managerial insights from these models, a task that is acknowledgedly difficult to pursue due to the complexity of the models;
- the other is the recognition of the gap between these tools and the tools that are required in practice to support corporate decisions, and the resulting need to set up research and development initiatives to bridge that gap.

On the other hand, concerning the application of multiobjective metaheuristics to stochastic combinatorial optimisation problems, we intend to focus on some incremental but natural developments to our work. Two possibilities of enhancing the Multiobjective Local Search template would be worth exploring:

- one would be to accommodate specific characteristics and components of stochastic problems, such as scenario trees or objective functions that are computed from a common vector of objective function values for each scenario;
- the other would be to generalise and incorporate the algorithmic solutions that have typically been used to adapt single-objective metaheuristics to stochastic optimisation - the incorporation of sampling methods for solution evaluation, and statistical inference methods for solution comparison.

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